# Detecting the Linkage in an $n$-Component Brunnian Link 

## IMUS Mini-Course Session 2



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Presented by Dr. Ron Umble
Millersville $U$ and IMUS
2 May 2018

## Recap of Lecture 1

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- Geometric diagonal $\Delta_{X}: X \rightarrow X \times X$
- Diagonal (approximation) $\Delta: C(X) \rightarrow C(X) \otimes C(X)$
- $\Delta$ induces a diagonal $\Delta_{2}: H(X) \rightarrow H(X) \otimes H(X)$ defined by

$$
\Delta_{2}[x]=[\Delta(x)]
$$

## Main result

$B R_{n}$ denote the link complement of an $n$-component Brunnian Link in $S^{3}$
Non-primitivity of the induced $\Delta_{2}$ on $H\left(B R_{2}\right)$ detects the Hopf link

## Goal of the Project

Use a similar strategy to detect linkage in an n-component Brunnian link


Borromean rings ( $n=3$ )

## The Borromean Rings



## A Tubular Neighborhood of the Borromean Rings



## A Cellular Decomposition of BR(3)



2-cells (front side view)

## A Cellular Decomposition of BR(3)



## A Cellular Decomposition of BR(3)



2-cells (top view)

## A Cellular Decomposition of BR(3)



Edges and vertices (front side view)

## A Cellular Decomposition of BR(3)



Edges and vertices (top view)

## Cellular Chains of BR(3)

- 11 vertices, 32 edges, 26 polygons, 5 solids

$$
\begin{aligned}
& C_{0}\left(B R_{3}\right)=\left\langle v_{1}, v_{1}, \ldots, v_{11}\right\rangle \\
& C_{1}\left(B R_{3}\right)=\left\langle m_{1}, \ldots, m_{14}, c_{1}, \ldots, c_{18}\right\rangle \\
& C_{2}\left(B R_{3}\right)=\left\langle a_{1}, \ldots, a_{4}, e_{1}, e_{2}, s_{1}, \ldots, s_{12}, t_{1}, \ldots, t_{8}\right\rangle \\
& C_{3}\left(B R_{3}\right)=\left\langle p, q_{1}, \ldots, q_{4}\right\rangle
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\end{aligned}
$$

- Euler characteristic $11-32+26-5=0$


## Cellular Boundary Map

- $\partial p=a_{1}+a_{2}+a_{3}+a_{4}$

$$
\begin{aligned}
& \partial q_{1}=a_{1}+e_{1}+s_{3}+s_{7}+s_{9}+t_{1}+t_{5} \\
& \partial q_{2}=a_{2}+e_{2}+s_{4}+s_{8}+s_{10}+t_{3}+t_{7} \\
& \partial q_{3}=a_{3}+e_{1}+s_{1}+s_{5}+s_{11}+t_{2}+t_{6} \\
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$$

- Boundaries of lower dim'l cells are evident from the pictures


## Cellular Homology of BR(3)

- $H_{0}\left(B R_{3}\right)=\left\langle\left[v_{1}\right]\right\rangle$
$H_{1}\left(B R_{3}\right)=\left\langle\left[m_{4}+m_{11}\right],\left[m_{7}+m_{8}\right],\left[c_{13}+c_{14}\right]\right\rangle$
$H_{2}\left(B R_{3}\right)=\left\langle\left[t_{1}+t_{2}+t_{3}+t_{4}\right],\left[t_{5}+t_{6}+t_{7}+t_{8}\right]\right\rangle$
$H_{k}\left(B R_{3}\right)=0, k \geq 3$


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$H_{2}\left(B R_{3}\right)=\left\langle\left[t_{1}+t_{2}+t_{3}+t_{4}\right],\left[t_{5}+t_{6}+t_{7}+t_{8}\right]\right\rangle$
$H_{k}\left(B R_{3}\right)=0, k \geq 3$
- Euler characteristic $1-3+2=0$


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- $\Delta m_{i}=\left(\right.$ minimal vertex of $\left.m_{i}\right) \otimes m_{i}+m_{i} \otimes\left(\right.$ maximal vertex of $\left.m_{i}\right)$


## Kravatz's Diagonal on a Polygon

Theorem (Kravatz, 2006) Let $G$ be an $n$-gon with initial vertex $v_{1}$, terminal vertex $v_{t}$, and edges $e_{1}, e_{1}, \ldots, e_{n}$ directed from $v_{1}$ to $v_{t}$. Then

$$
\Delta(G)=v_{1} \otimes G+G \otimes v_{t}+\sum_{0<i<j<t} e_{i} \otimes e_{j}+\sum_{n \geq j>i \geq t} e_{j} \otimes e_{i}
$$

defines a diagonal on $C(G)$.


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- $\Delta a_{i}=v_{1} \otimes a_{i}+a_{i} \otimes v_{11}$
- $\Delta t_{1}=v_{1} \otimes t_{1}+t_{1} \otimes v_{5}+m_{11} \otimes\left(c_{3}+c_{4}+m_{10}\right)$

$$
+c_{3} \otimes\left(c_{4}+m_{10}\right)+c_{4} \otimes m_{10}
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$$
\begin{aligned}
& +\left(c_{1}+c_{15}\right) \otimes t_{5}+t_{5} \otimes c_{13} \\
& +c_{1} \otimes s_{3}+s_{3} \otimes\left(c_{6}+m_{8}+c_{13}\right) \\
& +\left(m_{11}+c_{3}\right) \otimes s_{7}+s_{7} \otimes\left(m_{8}+c_{13}\right)
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## Conjecture

An extension of $\Delta$ to an $A_{\infty}$-coalgebra structure on $C\left(B R_{n}\right)$ induces an $A_{\infty}$-coalgebra structure on $H\left(B R_{n}\right)$ with

- A primitive diagonal

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- Detects the linkage in a n-component Brunnian link


## Minnich's A-infinity Coalgebra Structure on a Polygon

Theorem (Minnich, 2017) Let $G$ be an n-gon with initial vertex $v_{1}$, terminal vertex $v_{t}$, and edges $e_{1}, e_{1}, \ldots, e_{n}$ directed from $v_{1}$ to $v_{t}$. Let $\Delta_{2}$ denote the Kravatz diagonal. For $k>2$ define

$$
\Delta_{k}(G)=\sum_{0<i_{1}<\cdots<i_{k}<t} e_{i_{1}} \otimes \cdots \otimes e_{i_{k}}+\sum_{n \geq i_{1}>\cdots>i_{k} \geq t} e_{i_{1}} \otimes \cdots \otimes e_{i_{k}} .
$$

Then $\left(C(G), \partial, \Delta_{2}^{\prime}, \Delta_{3}^{\prime}, \ldots\right)$ is an $A_{\infty}$-coalgebra.


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- $(V, \partial)$ is a differential graded vector space (d.g.v.s.)


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- $H\left(\operatorname{Hom}_{*}(V, W)\right)=\operatorname{ker} \delta / \operatorname{Im} \delta$


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- When $p=0$ we have

$$
\begin{array}{lllll}
\cdots \longleftarrow & V_{i-1} & \longleftarrow \partial_{V} & V_{i} & \longleftarrow \\
& T \searrow & \downarrow^{f+g} & \searrow T & \\
& \cdots & \longleftarrow & W_{i} & \longleftarrow \\
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- $f+g$ is a boundary if $\delta(T)=f+g$ for some chain homotopy $T$


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- The associahedron $K_{n}$ is an ( $n-2$ )-dimensional polytope that controls homotopy (co)associativity in $n$ variables

- Associahedra organize the structural data in the definition of an $A_{\infty}$-(co)algebra
- For each $n \geq 2$, let $\theta_{n}$ denote the ( $n-2$ )-dimensional cell of $K_{n}$


## A-infinity Coalgebras Defined

- Let $(V, \partial)$ be a d.g.v.s. For each $n \geq 2$, choose a map $\alpha_{n}$ of deg 0 :

$$
\begin{array}{ccc}
C_{*}\left(K_{n}\right) & \xrightarrow[\alpha_{n}]{ } & \operatorname{Hom}_{*}\left(V, V^{\otimes n}\right) \\
\partial \downarrow & & \downarrow \delta \\
C_{*-1}\left(K_{n}\right) & \xrightarrow[\alpha_{n}]{\longrightarrow} & \operatorname{Hom}_{*-1}\left(V, V^{\otimes n}\right)
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- $\left(V, \partial, \Delta_{2}, \Delta_{3}, \ldots\right)$ is an $A_{\infty}$-coalgebra if each $\alpha_{n}$ is a chain map, i.e.,

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- Evaluating at $\theta_{n}$ produces the classical structure relations

$$
\delta\left(\Delta_{n}\right)=\sum_{i=1}^{n-2} \sum_{j=0}^{n-i-1}(-1)^{i(n+j+1)}\left(\mathbf{1}^{\otimes j} \otimes \Delta_{i+1} \otimes \mathbf{1}^{\otimes n-i-j-1}\right) \Delta_{n-i}
$$

## Structure Relations

$\Delta_{n}$ is a chain homotopy among the quadratic compositions encoded by the codim 1 cells of $K_{n}$


$$
\left(\Delta_{3} \otimes 1\right) \Delta_{2}
$$

$$
\begin{aligned}
& \delta\left(\Delta_{4}\right)=(\partial \otimes \mathbf{1} \otimes \mathbf{1}+\mathbf{1} \otimes \partial \otimes \mathbf{1}+\mathbf{1} \otimes \mathbf{1} \otimes \partial) \Delta_{4}+\Delta_{4} \partial \\
& \quad=\left(\Delta_{2} \otimes \mathbf{1} \otimes \mathbf{1}+\mathbf{1} \otimes \Delta_{2} \otimes \mathbf{1}+\mathbf{1} \otimes \mathbf{1} \otimes \Delta_{2}\right) \Delta_{3}+\left(\Delta_{3} \otimes \mathbf{1}+\mathbf{1} \otimes \Delta_{3}\right) \Delta_{2}
\end{aligned}
$$

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- $\Delta_{4}\left(t_{1}\right)=m_{11} \otimes c_{3} \otimes c_{4} \otimes m_{10}$
- $\Delta_{5}\left(s_{9}\right)=m_{11} \otimes c_{3} \otimes m_{13} \otimes m_{8} \otimes c_{13}$
- $\Delta_{k}=0$ on 2-cells for all $k \geq 6$


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- $\Delta_{3}\left(q_{1}\right)=t_{1} \otimes m_{12} \otimes\left(c_{6}+m_{8}+c_{13}\right)$

$$
\begin{aligned}
& +t_{1} \otimes c_{6} \otimes\left(m_{8}+c_{13}\right)+t_{1} \otimes m_{8} \otimes c_{13} \\
& +s_{3} \otimes c_{6} \otimes\left(m_{8}+c_{13}\right)+s_{3} \otimes m_{8} \otimes c_{13} \\
& +s_{7} \otimes m_{8}+c_{13}+\left(c_{1}+c_{15}\right) \otimes t_{5} \otimes c_{13} \\
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## Operations on 3-cells of $\operatorname{BR}(3)$

- M. Fansler (2016) computed $\Delta_{3}$ on 3-cells, e.g.,
- $\Delta_{3}\left(q_{1}\right)=t_{1} \otimes m_{12} \otimes\left(c_{6}+m_{8}+c_{13}\right)$

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- $\Delta_{k}=0$ for all $k \geq 6$


## Transferring Coproducts

## Goal:

$A_{\infty}$-coalgebra on chains
$\left(C, \partial, \Delta_{2}, \Delta_{3}, \ldots\right)$
$\downarrow$
$\left(H, 0, \Delta^{2}, \Delta^{3}, \ldots\right)$
$A_{\infty}$-coalgebra in homology

## Transferring Coproducts

Required input:

- Coalgebra on chains ( $C, \partial, \Delta_{2}, \Delta_{3}, \ldots$ ) and
- a cycle-selecting map $g: H \rightarrow Z(C)$, where $Z(C)$ denotes the subspace of cycles in $C$.
Note: In practice we only required $\Delta_{2}$ at the outset and computed the rest as needed.


## How Does It Work?

Strategy: Construct a chain map from the top dimension and codim-1 cells of the ( $n-1$ )-dimensional multiplihedron, denoted $J_{n}$, to maps between $H$ and $C^{\otimes n}$.

## Beginning Steps

- $J_{n}$ is a polytope that captures the combinatiorial structure of mapping between two $A_{\infty}$-coalgebras.


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- Consider $J_{1}$ and $J_{2}$.


## Extending to $J_{3}$



## Table of Contents

(1) Introduction
(2) Transfer Algorithm
(3) Implementation

4 Examples
(5) Conclusions

## Linear Algebraic Methods

## Good News

Linear algebra provides robust and theoretically correct methods for solving the various induction steps of the transfer algorithm.

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Linear algebra provides robust and theoretically correct methods for solving the various induction steps of the transfer algorithm.

## Bad News

The matrices are too large to be solved within a reasonable amount of storage space and time.

## Two Problems

## Problem (Preboundary)

Given a cycle $x \in C^{\otimes n}$ of degree $k$, find a chain $y \in C^{\otimes n}$ of degree $k+1$, such that $\partial(y)=x$.

## Problem (Factorization)

Given a cycle $c \in Z\left(C^{\otimes n}\right)$, find all subcycles of $c$ of the form $Z(C)^{\otimes n}$.

## Preboundary Problem: $\Delta_{3}$

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- Instead, solved with a best-first search algorithm


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- Again, an algorithmic approach appears to be a feasible alternative


## Induced Operations Computed by M. Fansler

- $H_{0}=\left\{0_{0}\right\}, H_{1}=\left\{1_{0}, 1_{1}, 1_{2}\right\}, H_{2}=\left\{2_{0}, 2_{1}\right\}$


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- Linkage detected but $\Delta^{4}$ remains to be computed


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- Her construction adjusts the decomposition of $B R_{3}$ so that all 2-cells have 5 edges
- Numbers of vertices, edges, faces, and solids in her decomposition are the same as in mine


## Nimershiem's Decomposition of BR(3)



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## The Case of BR(n)

- Redo the $B R_{3}$ calculations using Nimershiem's decomposition


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- Use Nimershiem's decomposition to calculate $B R_{4}$
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- Stay tuned!!


## The End

## Thank you!

