

## Order Relations

A **partial order** on a set is, roughly speaking, a relation that behaves like the relation  $\leq$  on  $\mathbb{R}$ .

**Definition.** Let  $X$  be a set, and let  $\sim$  be a relation on  $X$ .  $\sim$  is a **partial order** if:

- (a) (Reflexive) For all  $x \in X$ ,  $x \sim x$ .
- (b) (Antisymmetric) For all  $x, y \in X$ , if  $x \sim y$  and  $y \sim x$ , then  $x = y$ .
- (c) (Transitive) For all  $x, y, z \in X$ , if  $x \sim y$  and  $y \sim z$ , then  $x \sim z$ .

**Example.** For each relation, check each axiom for a partial order. If the axiom holds, prove it. If the axiom does not hold, give a specific counterexample.

(a) The relation  $\leq$  on  $\mathbb{R}$ .

(b) The relation  $<$  on  $\mathbb{R}$ .

(a) For all  $x \in \mathbb{R}$ ,  $x \leq x$ : Reflexivity holds.

For all  $x, y \in \mathbb{R}$ , if  $x \leq y$  and  $y \leq x$ , then  $x = y$ : Antisymmetry holds.

For all  $x, y, z \in \mathbb{R}$ , if  $x \leq y$  and  $y \leq z$ , then  $x \leq z$ : Transitivity holds.

Thus,  $\leq$  is a partial order.  $\square$

(b) For no  $x$  is it true that  $x < x$ , so reflexivity fails.

Antisymmetry would say: If  $x < y$  and  $y < x$ , then  $x = y$ . However, for no  $x, y \in \mathbb{R}$  is it true that  $x < y$  and  $y < x$ . Therefore, the first part of the conditional is false, and the conditional is true. Thus, antisymmetry is *vacuously true*.

If  $x < y$  and  $y < z$ , then  $x < z$ . Therefore, transitivity holds.

Hence,  $<$  is not a partial order.  $\square$

**Example.** Let  $X$  be a set and let  $\mathcal{P}(X)$  be the power set of  $X$  — i.e. the set of all subsets of  $X$ . Show that the relation of **set inclusion** is a partial order on  $\mathcal{P}(X)$ .

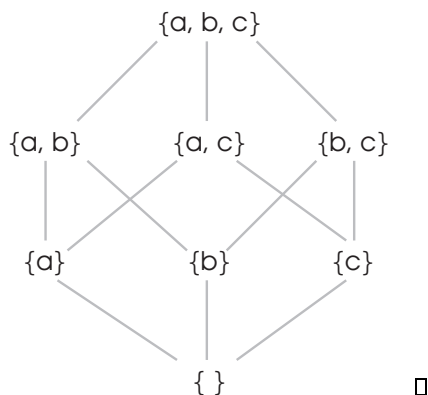
Subsets  $A$  and  $B$  of  $X$  are related under set inclusion if  $A \subset B$ .

If  $A \subset X$ , then  $A \subset A$ . The relation is reflexive.

Suppose  $A, B \subset X$ . If  $A \subset B$  and  $B \subset A$ , then by definition of set equality,  $A = B$ . The relation is symmetric.

Finally, suppose  $A, B, C \subset X$ . If  $A \subset B$  and  $B \subset C$ , then  $A \subset C$ . (You can write out the easy proof using elements.) The relation is transitive.

Here's a particular example. Let  $X = \{a, b, c\}$ . This is a picture of the set inclusion relation on  $\mathcal{P}(X)$ :



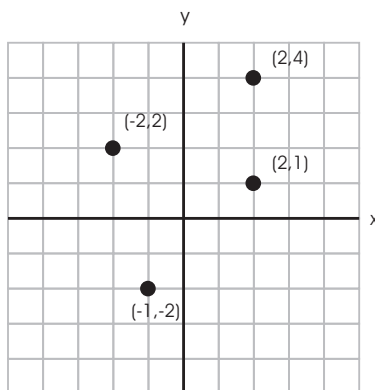
□

**Definition.** Let  $(X, \leq)$  be a partially ordered set. The **lexicographic order** (or **dictionary order**) on  $X \times X$  is defined as follows:  $(x_1, y_1) \sim (x_2, y_2)$  means that

- (a)  $x_1 < x_2$ , or
- (b)  $x_1 = x_2$  and  $y_1 \leq y_2$ .

Note that  $(x_1, y_1) \sim (x_2, y_2)$  implies  $x_1 \leq x_2$ .

You can extend the definition to two different partially ordered sets  $X$  and  $Y$ , or a sequence  $X_1, X_2, \dots, X_n$  of partially ordered sets in the same way. The name *dictionary order* comes from the fact that it describes the way words are ordered alphabetically in a dictionary. For instance, “aardvark” comes before “banana” because “a” comes before “b”. If the first letters are the same, as with “mystery” and “meat”, then you look at the second letters: “e” comes before “y”, so “meat” comes before “mystery”.



In the picture above,  $(-2, 2) \sim (-1, -2)$ , because  $-2 < -1$ . And  $(2, 1) \sim (2, 4)$  because the  $x$ -coordinates are equal and  $1 < 4$ .

**Proposition.** The lexicographic order on  $X \times X$  is a partial order.

**Proof.** First,  $(x, y) \sim (x, y)$ , since  $x = x$  and  $y \leq y$ .  $\sim$  is reflexive.

Next, suppose  $(x_1, y_1) \sim (x_2, y_2)$  and  $(x_2, y_2) \sim (x_1, y_1)$ . Now  $(x_1, y_1) \sim (x_2, y_2)$  means that either  $x_1 < x_2$  or  $x_1 = x_2$ . The first case  $x_1 < x_2$  is impossible, since this would contradict  $(x_2, y_2) \sim (x_1, y_1)$ . Therefore,  $x_1 = x_2$ . Then  $(x_1, y_1) \sim (x_2, y_2)$  implies  $y_1 \leq y_2$  and  $(x_2, y_2) \sim (x_1, y_1)$  implies  $y_2 \leq y_1$ . Hence,  $y_1 = y_2$ . Therefore,  $(x_1, y_1) = (x_2, y_2)$ .  $\sim$  is antisymmetric.

Finally, suppose  $(x_1, y_1) \sim (x_2, y_2)$  and  $(x_2, y_2) \sim (x_3, y_3)$ . To keep things organized, I'll consider the four cases.

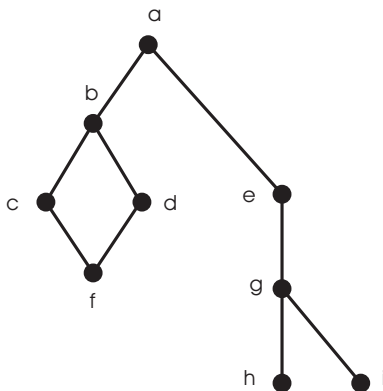
- (a) If  $x_1 < x_2$  and  $x_2 < x_3$ , then  $x_1 < x_3$ , so  $(x_1, y_1) \sim (x_3, y_3)$ .
- (b) If  $x_1 < x_2$  and  $x_2 = x_3$ , then  $x_1 < x_3$ , so  $(x_1, y_1) \sim (x_3, y_3)$ .
- (c) If  $x_1 = x_2$  and  $x_2 < x_3$ , then  $x_1 < x_3$ , so  $(x_1, y_1) \sim (x_3, y_3)$ .
- (d) If  $x_1 = x_2$  and  $x_2 = x_3$ , then  $y_1 \leq y_2$  and  $y_2 \leq y_3$ . This implies  $x_1 = x_3$  and  $y_1 \leq y_3$ , so  $(x_1, y_1) \sim (x_3, y_3)$ .

Hence,  $\sim$  is transitive, and this completes the proof that  $\sim$  is a partial order.  $\square$

A common mistake in working with partial orders — and in real life — consists of *assuming* that if you have two things, then one must be bigger than the other. When this *is* true about two things, the things are said to be **comparable**. However, in an arbitrary partially ordered set, some pairs of elements are comparable and some are not.

**Definition.** Let  $\sim$  be a relation on a set  $X$ .  $x$  and  $y$  in  $X$  are **comparable** if either  $x \sim y$  or  $y \sim x$ .

Here’s a pictorial example to illustrate the idea. You can sometimes describe an order relation by drawing a graph like the one below:



This picture shows a relation  $\sim$  on the set

$$S = \{a, b, c, d, e, f, g, h, i\}.$$

Two elements are comparable if they’re joining by a sequence of edges that goes upward “without reversing direction”. (Think of “bigger” elements being above and “smaller” elements being below.) It’s also understood that every element satisfies  $x \sim x$ .

For example,  $f \sim c$ , since there’s an upward segment connecting  $f$  to  $c$ . And  $f \sim a$ , since there’s an upward path of segments  $f \rightarrow c \rightarrow b \rightarrow a$  connecting  $f$  to  $a$ .

On the other hand, there are elements which are not comparable. For example,  $d$  and  $e$  are not comparable, because there is no upward path of segments connecting one to the other. Likewise,  $g \sim h$  and  $g \sim i$ , but  $h$  and  $i$  are not comparable.

Notice that  $a$  is comparable to every element of the set, and that  $x \sim a$  for all  $x \in S$ .

**Definition.** Let  $X$  be a partially ordered set.

- (a) An element  $x \in X$  which is comparable to every other element of  $X$  and satisfies  $x \geq y$  for all  $y \in X$  is the **largest element** of the set.
- (b) An element  $x \in X$  which is comparable to every other element of  $X$  and satisfies  $x \leq y$  for all  $y \in X$  is the **smallest element** of the set.

In some cases, we only care that an element be “bigger than” or “smaller than” elements to which it is comparable.

**Definition.** Let  $X$  be a partially ordered set. If an element  $x$  satisfies  $x \geq y$  for all  $y$  to which it is comparable, then  $x$  is a **maximal element**. Likewise, if an element  $x$  satisfies  $x \leq y$  for all  $y$  to which it is comparable, then  $x$  is a **minimal element**.

Note that a largest or smallest element, if it exists, is unique. On the other hand, there may be many maximal or minimal elements.

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**Example.** Define a relation  $\sim$  on  $\mathbb{R}$  by

$$x \sim y \quad \text{means} \quad x^3 - 4x \leq y^3 - 4y.$$

Check each axiom for a partial order. If the axiom holds, prove it. If the axiom does not hold, give a specific counterexample.

$x^3 - 4x \leq x^3 - 4x$  for all  $x \in \mathbb{R}$ , so  $x \sim x$  for all  $x \in \mathbb{R}$ . Therefore,  $\sim$  is reflexive.

Suppose  $x \sim y$  and  $y \sim x$ . Is it true that  $x = y$ ?

$2 \sim -2$ , since  $2^3 - 4 \cdot 2 \leq (-2)^3 - 4 \cdot (-2)$ . Likewise,  $-2 \sim 2$ , since  $(-2)^3 - 4 \cdot (-2) \leq 2^3 - 4 \cdot 2$ . But  $2 \neq -2$ , so  $\sim$  is not antisymmetric.

Finally, suppose  $x \sim y$  and  $y \sim z$ . This means that  $x^3 - 4x \leq y^3 - 4y$  and  $y^3 - 4y \leq z^3 - 4z$ . Hence,  $x^3 - 4x \leq z^3 - 4z$ . Therefore,  $x \sim z$ , so  $\sim$  is transitive.  $\square$

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**Example.** Define a relation  $\sim$  on  $\mathbb{R}^2$  by

$$(a, b) \sim (c, d) \quad \text{means} \quad |ab| \geq |cd|.$$

Check each axiom for a partial order. If the axiom holds, prove it. If the axiom does not hold, give a specific counterexample.

Since  $|ab| \geq |ab|$  for all  $(a, b) \in \mathbb{R}^2$ , it follows that  $(a, b) \sim (a, b)$  for all  $(a, b) \in \mathbb{R}^2$ . Therefore,  $\sim$  is reflexive.

$(1, 2) \sim (-1, 2)$ , since  $|1 \cdot 2| \geq |(-1) \cdot 2|$ . Likewise,  $(-1, 2) \sim (1, 2)$ , since  $|(-1) \cdot 2| \geq |1 \cdot 2|$ . However,  $(1, 2) \not\sim (-1, 2)$ . Therefore,  $\sim$  is not antisymmetric.

Finally, suppose  $(a, b) \sim (c, d)$  and  $(c, d) \sim (e, f)$ . Then  $|ab| \geq |cd|$  and  $|cd| \geq |ef|$ . Hence,  $|ab| \geq |ef|$ . Therefore,  $(a, b) \sim (e, f)$ . Hence,  $\sim$  is transitive.  $\square$

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**Definition.** A relation  $\sim$  on a set  $X$  is a **total order** if:

- (a) (Trichotomy) For all  $x, y \in X$ , exactly one of the following holds:  $x \sim y$ ,  $y \sim x$ , or  $x = y$ .
- (b) (Transitivity) For all  $x, y, z \in X$ , if  $x \sim y$  and  $y \sim z$ , then  $x \sim z$ .

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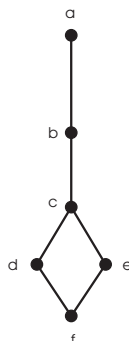
The usual less than relation  $<$  is a total order on  $\mathbb{Z}$ , on  $\mathbb{Q}$ , and on  $\mathbb{R}$ . Likewise, you can use the total order relation on  $\mathbb{Z}$  to define a lexicographic order on  $\mathbb{Z} \times \mathbb{Z}$  which is a total order. Specifically, define a total order  $\sim$  on  $\mathbb{Z} \times \mathbb{Z}$  as follows:  $(x_1, y_1) \sim (x_2, y_2)$  means that

- (a)  $x_1 < x_2$ , or

(b)  $x_1 = x_2$  and  $y_1 < y_2$ .

You can check that the axioms for a total order hold.

**Example.** Consider the relation defined by the graph below:



Thus,  $x < y$  means that  $x \neq y$ , and there is an upward path of segments from  $x$  to  $y$ .

Is this relation a total order? You can check cases, using the picture, that the relation is transitive. (This amounts to saying that if there's an upward path from  $x$  to  $y$  and one from  $y$  to  $z$ , then there's such a path from  $x$  to  $z$ . In fact, if you define a relation using a graph in this way, the relation will be transitive.)

However, this graph does not define a total order. Trichotomy fails for  $d$  and  $e$ , since  $d < e$ ,  $e < d$ , and  $d = e$  are *all* false.  $\square$

**Definition.** Let  $S$  be a partially ordered set, and let  $T$  be a subset of  $S$ .

(a)  $s \in S$  is an **upper bound** for  $T$  if  $s \geq t$  for all  $t \in T$ .

(b)  $s \in S$  is a **lower bound** for  $T$  if  $s \leq t$  for all  $t \in T$ .

Thus, an upper bound for a subset is an element which is greater than or equal to everything in the subset; a lower bound for a subset is an element which is less than or equal to everything in the subset. Note that unlike the **largest element** or **smallest element** of a subset, upper and lower bounds don't need to belong to the subset.

For instance, consider the subset  $T = (0, 1]$  of  $\mathbb{R}$ . 2 is an upper bound for  $T$ , since  $2 \geq x$  for all  $x \in T$ . 1 is also an upper bound for  $T$ . Note that 2 is not an element of  $T$  while 1 is an element of  $T$ . In fact, any real number greater than or equal to 1 is an upper bound for  $T$ .

Likewise, any real number less than or equal to 0 is a lower bound for  $T$ .

$T$  has a largest element, namely 1. It does not have a smallest element; the obvious candidate 0 is not in  $T$ .

This example shows that a subset may have many — even infinitely many — upper or lower bounds. Among all the upper bounds for a set, there may be one which is *smallest*.

**Definition.** Let  $S$  be a partially ordered set, and let  $T$  be a subset of  $S$ . An element  $s_0 \in S$  is a **least upper bound** for  $T$  if:

(a)  $s_0$  is an upper bound for  $T$ .

(b) If  $s$  is an upper bound for  $T$ , then  $s_0 \leq s$ .

The idea is that  $s_0$  is an upper bound by (a); it's the **least** upper bound, since (b) says  $s_0$  is smaller than any other upper bound.

**Definition.** Let  $S$  be a partially ordered set, and let  $T$  be a subset of  $S$ . An element  $s_0 \in S$  is a **greatest lower bound** for  $T$  if:

- (a)  $s_0$  is an lower bound for  $T$ .
- (b) If  $s$  is an lower bound for  $T$ , then  $s_0 \geq s$ .

The concepts of least upper bound and greatest lower bound come up often in analysis. I'll give a simple example.

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**Example.** Determine the least upper bound and greatest lower bound for the following sets (if they exist):

- (a) The subset  $S = (0, 1]$  of  $\mathbb{R}$ .
- (b) The subset  $T = (0, +\infty)$  of  $\mathbb{R}$ . (Thus,  $T$  is the positive real axis, not including 0.)

(a) Any real number greater than or equal to 1 is an upper bound for  $T$ . Among the upper bounds for  $S$ , it's clear that 1 is the *smallest*, so 1 is the *least upper bound* for  $S$ .

Likewise, any real number less than or equal to 0 is a lower bound for  $S$ . But among the lower bounds for  $S$ , it's clear that 0 is the *largest*, so 0 is the *greatest lower bound* for  $S$ .

Notice that  $1 \in S$ , but  $0 \notin S$ . The least upper bound and greatest lower bound may be contained, or not contained, in the set.  $\square$

- (b)  $T$  has no least upper bound in  $\mathbb{R}$ ; in fact,  $T$  has no upper bound in  $\mathbb{R}$ .  
0 is the greatest lower bound for  $T$  in  $\mathbb{R}$ .  $\square$
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