

Non-operadic Operations on Loop Cohomology

Joint work with Samson Saneblidze

Ron Umble

Millersville Univ of Pennsylvania

Modern Algebra and its Applications

September 19, 2011

Goal of the talk

- To identify a space X such that $H = H^*(\Omega X; \mathbb{Z}_2)$ comes equipped with an induced operation

$$\omega_2^2 : H \otimes H \rightarrow H \otimes H$$

Goal of the talk

- To identify a space X such that $H = H^*(\Omega X; \mathbb{Z}_2)$ comes equipped with an induced operation

$$\omega_2^2 : H \otimes H \rightarrow H \otimes H$$

- We assume existence of biassociahedra $KK = \{KK_{n,m}\}$ and bmultipliheda $JJ = \{JJ_{n,m}\}$

Goal of the talk

- To identify a space X such that $H = H^*(\Omega X; \mathbb{Z}_2)$ comes equipped with an induced operation

$$\omega_2^2 : H \otimes H \rightarrow H \otimes H$$

- We assume existence of biassociahedra $KK = \{KK_{n,m}\}$ and bmultipliheda $JJ = \{JJ_{n,m}\}$
- The operation ω_2^2 will be defined in terms of $JJ_{2,1} = JJ_{1,2}$ and $JJ_{2,2}$

The Free Matrad

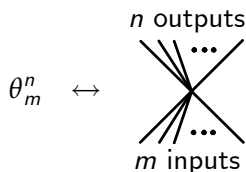
- $\mathcal{H}_\infty$ is matradically generated by $\{\theta_m^n \mid mn \neq 1\}$

The Free Matrad

- $\mathcal{H}_\infty$ is matradically generated by $\{\theta_m^n \mid mn \neq 1\}$
- $\mathcal{H}_\infty$ is realized by cellular chains $C_*(KK)$

The Free Matrad

- $\mathcal{H}_\infty$ is matratically generated by $\{\theta_m^n \mid mn \neq 1\}$
- $\mathcal{H}_\infty$ is realized by cellular chains $C_*(KK)$
- Top dim'l cell of $KK_{n,m}$ is identified with the matradic generator



The Free Matrad

- Cells of KK are identified with certain fraction product monomials

The diagram shows an equality between a fraction of two monomials and a vertex of $KK_{2,3}$. On the left, a fraction has a numerator consisting of two separate trivalent vertices (each with one top and two bottom legs) and a denominator consisting of three separate trivalent vertices (each with two top and one bottom leg). This fraction is equal to a single complex diagram on the right, which is a vertex of $KK_{2,3}$. This vertex is represented as a diamond shape with internal lines forming a more complex trivalent structure. To the right of this diagram is a double-headed arrow pointing to the text "a vertex of $KK_{2,3}$ ".

The Free Matrad

- Cells of KK are identified with certain fraction product monomials

$$\frac{\begin{array}{cc} \text{Y-shape} & \text{Y-shape} \\ \hline \text{Y-shape} & \text{Y-shape} & \text{Y-shape} \end{array}}{=} \begin{array}{c} \text{A complex diagram consisting of two diamond shapes sharing a vertex, with internal lines and a dashed line.} \end{array} \leftrightarrow \text{a vertex of } KK_{2,3}$$

- $KK_{n,1} = KK_{1,n} = K_n$ (Stasheff's associahedron)

The Free Matrad

- Cells of KK are identified with certain fraction product monomials

$$\frac{\begin{array}{cc} \text{Y-shape} & \text{Y-shape} \\ \hline \text{Y-shape} & \text{Y-shape} & \text{Y-shape} \end{array}}{=} \begin{array}{c} \text{A complex diagram consisting of two diamond shapes sharing a vertex, with internal lines and a dashed line.} \end{array} \leftrightarrow \text{a vertex of } KK_{2,3}$$

- $KK_{n,1} = KK_{1,n} = K_n$ (Stasheff's associahedron)
- $\dim(KK_{n,m}) = m + n - 3$

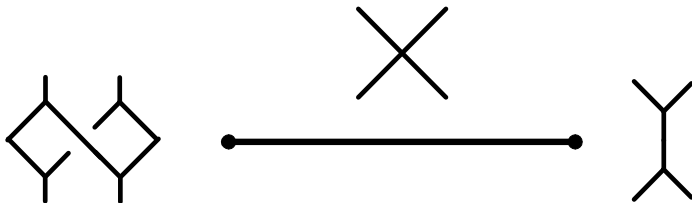
The Free Matrad

- Cells of KK are identified with certain fraction product monomials

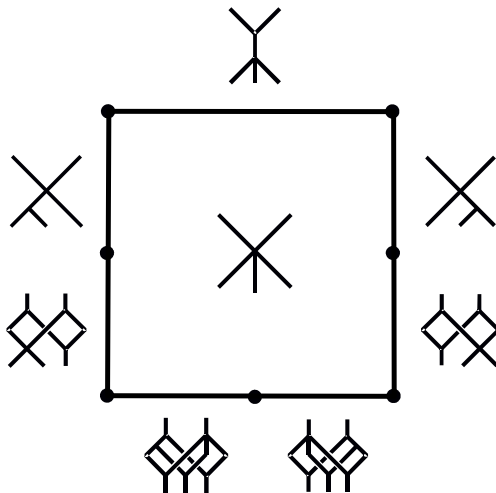
$$\frac{\begin{array}{cc} \text{Y-shape} & \text{Y-shape} \\ \hline \text{Y-shape} & \text{Y-shape} & \text{Y-shape} \end{array}}{=} \begin{array}{c} \text{A complex graph with 6 vertices and 9 edges, representing a vertex of } KK_{2,3} \end{array} \leftrightarrow \text{a vertex of } KK_{2,3}$$

- $KK_{n,1} = KK_{1,n} = K_n$ (Stasheff's associahedron)
- $\dim(KK_{n,m}) = m + n - 3$
- $KK_{n,m}$ is contractible

The Biassociahedron $KK(2,2)$:



The Biassociahedron $KK(2,3)$:



- Let (A, d) be a DGM over a commutative ring R with unity

A-infinity Bialgebras

- Let (A, d) be a DGM over a commutative ring R with unity
- $\mathcal{E}nd_{TA}$ is a matrad with respect to $\odot$ -product

A-infinity Bialgebras

- Let (A, d) be a DGM over a commutative ring R with unity
- $\mathcal{E}nd_{TA}$ is a matrad with respect to $\odot$ -product
- (A, d) together with a family of multilinear operations

$$\omega = \{ \omega_m^n \in \text{Hom}(A^{\otimes m}, A^{\otimes n}) \mid mn > 1 \}$$

is an A_∞ -bialgebra iff either

A-infinity Bialgebras

- Let (A, d) be a DGM over a commutative ring R with unity
- $\mathcal{E}nd_{TA}$ is a matrad with respect to $\odot$ -product
- (A, d) together with a family of multilinear operations

$$\omega = \{ \omega_m^n \in \text{Hom}(A^{\otimes m}, A^{\otimes n}) \mid mn > 1 \}$$

is an A_∞ -bialgebra iff either

- $\omega \odot \omega = 0$ or

A-infinity Bialgebras

- Let (A, d) be a DGM over a commutative ring R with unity
- $\mathcal{E}nd_{TA}$ is a matrad with respect to $\odot$ -product
- (A, d) together with a family of multilinear operations

$$\omega = \{ \omega_m^n \in \text{Hom}(A^{\otimes m}, A^{\otimes n}) \mid mn > 1 \}$$

is an A_∞ -bialgebra iff either

- $\omega \odot \omega = 0$ or
- ω is the image of map $\mathcal{H}_\infty \rightarrow \mathcal{E}nd_{TA}$ of matrads sending $\theta_m^n \mapsto \omega_m^n$

- Let (A, d) be a DGM over a commutative ring R with unity
- $\mathcal{E}nd_{TA}$ is a matrad with respect to $\odot$ -product
- (A, d) together with a family of multilinear operations

$$\omega = \{ \omega_m^n \in \text{Hom}(A^{\otimes m}, A^{\otimes n}) \mid mn > 1 \}$$

is an A_∞ -bialgebra iff either

- $\omega \odot \omega = 0$ or
- ω is the image of map $\mathcal{H}_\infty \rightarrow \mathcal{E}nd_{TA}$ of matrads sending $\theta_m^n \mapsto \omega_m^n$
- An A_∞ -bialgebra (A, d, ω) is an algebra over $\mathcal{H}_\infty$

The Free Relative Matrad

- $\mathcal{II}_\infty$ is an $\mathcal{H}_\infty$ -bimodule generated by $\{f_m^n\}$

The Free Relative Matrad

- $\mathcal{J}\mathcal{J}_\infty$ is an $\mathcal{H}_\infty$ -bimodule generated by $\{f_m^n\}$
- $\mathcal{J}\mathcal{J}_\infty$ is realized by cellular chains $C_*(JJ)$

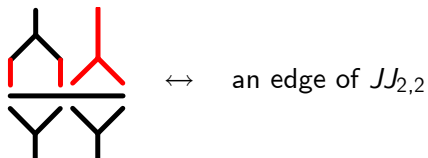
The Free Relative Matrad

- $\mathcal{JJ}_\infty$ is an $\mathcal{H}_\infty$ -bimodule generated by $\{\mathfrak{f}_m^n\}$
- $\mathcal{JJ}_\infty$ is realized by cellular chains $C_*(JJ)$
- Top dim'l cell of $JJ_{n,m}$ is identified with the $\mathcal{H}_\infty$ -bimodule generator

$$\mathfrak{f}_m^n \leftrightarrow \begin{array}{c} n \\ \diagup \quad \diagdown \\ \vdots \quad \vdots \\ \diagdown \quad \diagup \\ m \end{array}$$

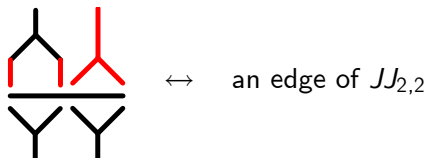
The Free Relative Matrad

- Cells of JJ are identified with certain fraction product monomials



The Free Relative Matrad

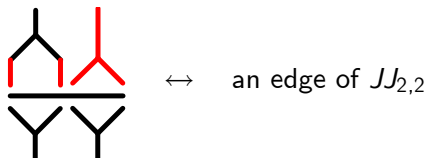
- Cells of JJ are identified with certain fraction product monomials



- $JJ_{n,1} = JJ_{1,n} = J_n$ (the multiplihedron)

The Free Relative Matrad

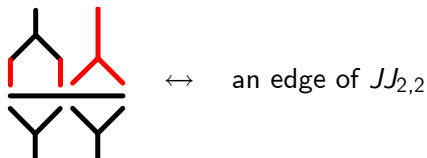
- Cells of JJ are identified with certain fraction product monomials



- $JJ_{n,1} = JJ_{1,n} = J_n$ (the multiplihedron)
- $\dim(JJ_{n,m}) = m + n - 2$

The Free Relative Matrad

- Cells of JJ are identified with certain fraction product monomials



- $JJ_{n,1} = JJ_{1,n} = J_n$ (the multiplihedron)
- $\dim(JJ_{n,m}) = m + n - 2$
- $JJ_{n,m}$ is a subdivision of $KK_{n,m} \times I$

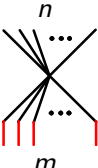
The Free Relative Matrad

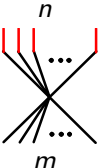
$KK_{n,m}$ embeds in $JJ_{n,m}$ as a codim 1 cell two ways:

$$\bullet \text{ As } KK_{n,m} \times 0 \iff \begin{array}{c} n \\ \diagup \quad \diagdown \\ \vdots \quad \vdots \\ \diagdown \quad \diagup \\ m \end{array} \iff \theta_m^n (f_1^1)^{\otimes m}$$

The Free Relative Matrad

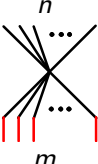
$KK_{n,m}$ embeds in $JJ_{n,m}$ as a codim 1 cell two ways:

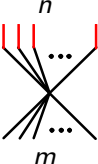
- As $KK_{n,m} \times 0 \leftrightarrow$  $\leftrightarrow \theta_m^n (f_1^1)^{\otimes m}$

- As $KK_{n,m} \times 1 \leftrightarrow$  $\leftrightarrow (f_1^1)^{\otimes n} \theta_m^n$

The Free Relative Matrad

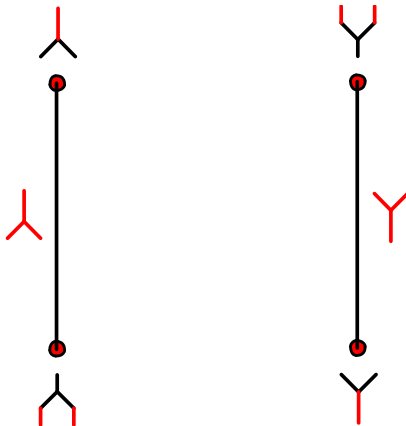
$KK_{n,m}$ embeds in $JJ_{n,m}$ as a codim 1 cell two ways:

- As $KK_{n,m} \times 0 \leftrightarrow$  $\leftrightarrow \theta_m^n (f_1^1)^{\otimes m}$

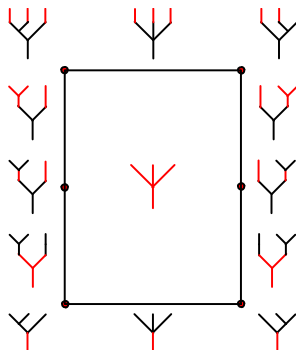
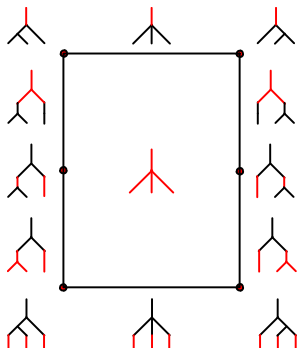
- As $KK_{n,m} \times 1 \leftrightarrow$  $\leftrightarrow (f_1^1)^{\otimes n} \theta_m^n$

- This latter cell plays an important role in our transfer algorithm

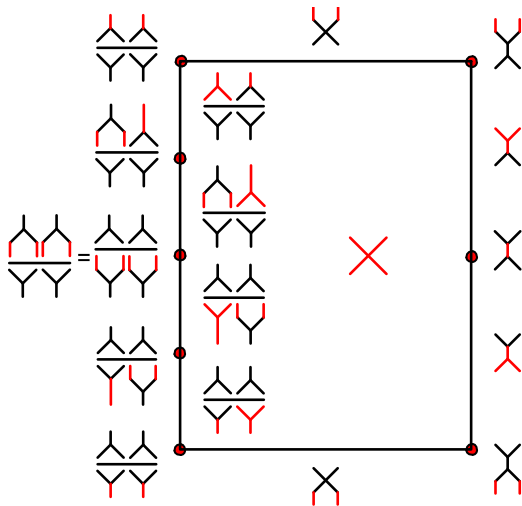
The Bimultiplihedra $JJ(1,2)$ and $JJ(2,1)$:



The Bimultiplihedra $JJ(1,3)$ and $JJ(3,1)$:



The Bimultiplihedron $JJ(2,2)$:



Morphisms of A-infinity Bialgebras

- Let (A, d_A, ω_A) and (B, d_B, ω_B) be A_∞ -bialgebras

Morphisms of A -infinity Bialgebras

- Let (A, d_A, ω_A) and (B, d_B, ω_B) be A_∞ -bialgebras
- $U_{A,B} = \text{Hom}(TA, TB)$ is a relative matrad
(as an $\mathcal{E}nd_{TB}$ - $\mathcal{E}nd_{TA}$ -bimodule)

Morphisms of A_∞ -bialgebras

- Let (A, d_A, ω_A) and (B, d_B, ω_B) be A_∞ -bialgebras
- $U_{A,B} = \text{Hom}(TA, TB)$ is a relative matrad
(as an $\mathcal{E}nd_{TB}$ - $\mathcal{E}nd_{TA}$ -bimodule)
- A morphism $A \Rightarrow B$ of A_∞ -bialgebras is the image

$$G = \{g_m^n \in \text{Hom}(A^{\otimes m}, B^{\otimes n})\}$$

of a map $\mathcal{J}\mathcal{J}_\infty \rightarrow U_{A,B}$ of relative matrads sending $f_m^n \mapsto g_m^n$

Morphisms of A_∞ -bialgebras

- Let (A, d_A, ω_A) and (B, d_B, ω_B) be A_∞ -bialgebras
- $U_{A,B} = \text{Hom}(TA, TB)$ is a relative matrad
(as an $\mathcal{E}nd_{TB}$ - $\mathcal{E}nd_{TA}$ -bimodule)

- A morphism $A \Rightarrow B$ of A_∞ -bialgebras is the image

$$G = \{g_m^n \in \text{Hom}(A^{\otimes m}, B^{\otimes n})\}$$

of a map $\mathcal{J}\mathcal{J}_\infty \rightarrow U_{A,B}$ of relative matrads sending $f_m^n \mapsto g_m^n$

- A morphism $G : A \Rightarrow B$ is an $\mathcal{H}_\infty$ -bimodule

Induced cochain map

- The differential ∇ on $f \in \text{Hom}(V, W)$ is given by

$$\nabla f = d_W f - (-1)^{|f|} f d_V$$

Induced cochain map

- The differential ∇ on $f \in \text{Hom}(V, W)$ is given by

$$\nabla f = d_W f - (-1)^{|f|} f d_V$$

- ∇ -cocycles are chain maps $V \rightarrow W$

Induced cochain map

- The differential ∇ on $f \in \text{Hom}(V, W)$ is given by

$$\nabla f = d_W f - (-1)^{|f|} f d_V$$

- ∇ -cocycles are chain maps $V \rightarrow W$
- A chain map $g : A \rightarrow B$ of DGMs induces a cochain map

$$\tilde{g} : \mathcal{E}nd_{TA} \rightarrow U_{A,B}$$

defined on $u \in \text{Hom}(A^{\otimes m}, A^{\otimes n})$ by $\tilde{g}(u) = g^{\otimes n} u$

The General Transfer Theorem – Generalizing the BPL

Let $g : A \rightarrow B$ be a chain map that is also a homology isomorphism

- **Proposition 1** *Then $\tilde{g}$ is a homology isomorphism if either*

The General Transfer Theorem – Generalizing the BPL

Let $g : A \rightarrow B$ be a chain map that is also a homology isomorphism

- **Proposition 1** *Then $\tilde{g}$ is a homology isomorphism if either*
 - *A is free as an R -module, or*

The General Transfer Theorem – Generalizing the BPL

Let $g : A \rightarrow B$ be a chain map that is also a homology isomorphism

- **Proposition 1** *Then $\tilde{g}$ is a homology isomorphism if either*
 - *A is free as an R -module, or*
 - *for each $n \geq 1$, there is a splitting $B^{\otimes n} = A^{\otimes n} \oplus X(n)$ such that $H^* \operatorname{Hom}(A^{\otimes k}, X(n)) = 0$ for all $k \geq 1$*

The General Transfer Theorem – Generalizing the BPL

Let $g : A \rightarrow B$ be a chain map that is also a homology isomorphism

- **Proposition 1** *Then $\tilde{g}$ is a homology isomorphism if either*
 - *A is free as an R -module, or*
 - *for each $n \geq 1$, there is a splitting $B^{\otimes n} = A^{\otimes n} \oplus X(n)$ such that $H^* \text{Hom}(A^{\otimes k}, X(n)) = 0$ for all $k \geq 1$*
- **Transfer Theorem** *If $\tilde{g}$ is a homology isomorphism, then*

The General Transfer Theorem – Generalizing the BPL

Let $g : A \rightarrow B$ be a chain map that is also a homology isomorphism

- **Proposition 1** *Then $\tilde{g}$ is a homology isomorphism if either*
 - *A is free as an R -module, or*
 - *for each $n \geq 1$, there is a splitting $B^{\otimes n} = A^{\otimes n} \oplus X(n)$ such that $H^* \text{Hom}(A^{\otimes k}, X(n)) = 0$ for all $k \geq 1$*
- **Transfer Theorem** *If $\tilde{g}$ is a homology isomorphism, then*
 - *g induces an A_∞ -bialgebra structure $\omega_A = \{\omega_A^{n,m}\}$ on A and*

The General Transfer Theorem – Generalizing the BPL

Let $g : A \rightarrow B$ be a chain map that is also a homology isomorphism

- **Proposition 1** *Then $\tilde{g}$ is a homology isomorphism if either*
 - *A is free as an R -module, or*
 - *for each $n \geq 1$, there is a splitting $B^{\otimes n} = A^{\otimes n} \oplus X(n)$ such that $H^* \text{Hom}(A^{\otimes k}, X(n)) = 0$ for all $k \geq 1$*
- **Transfer Theorem** *If $\tilde{g}$ is a homology isomorphism, then*
 - *g induces an A_∞ -bialgebra structure $\omega_A = \{\omega_A^{n,m}\}$ on A and*
 - *g extends to an A_∞ -bialgebra map $G = \{g_m^n \mid g_1^1 = g\} : A \Rightarrow B$*

The Transfer Algorithm: Given initial data

- A DGM (A, d_A)

The Transfer Algorithm: Given initial data

- A DGM (A, d_A)
- An A_∞ -bialgebra (B, d_B, ω_B) and a map of matrads

$$\alpha_B : C_*(KK) \rightarrow \mathcal{E}nd_{TB} \text{ sending } \theta_m^n \mapsto \omega_B^{n,m}$$

The Transfer Algorithm: Given initial data

- A DGM (A, d_A)
- An A_∞ -bialgebra (B, d_B, ω_B) and a map of matrads

$$\alpha_B : C_*(KK) \rightarrow \mathcal{E}nd_{TB} \text{ sending } \theta_m^n \mapsto \omega_B^{n,m}$$

- A chain map/homology isomorphism $g : A \rightarrow B$ such that $\tilde{g}$ is a homology isomorphism

The Transfer Algorithm: Objectives

- Define operations $\omega_A^{n,m} : A^{\otimes m} \rightarrow A^{\otimes n}$ for all $m, n, mn \neq 1$

The Transfer Algorithm: Objectives

- Define operations $\omega_A^{n,m} : A^{\otimes m} \rightarrow A^{\otimes n}$ for all $m, n, mn \neq 1$
- Define a map of matrads $\alpha_A : C_*(KK) \rightarrow \mathcal{E}nd_{TA}$ sending $\theta_m^n \mapsto \omega_A^{n,m}$

The Transfer Algorithm: Objectives

- Define operations $\omega_A^{n,m} : A^{\otimes m} \rightarrow A^{\otimes n}$ for all $m, n, mn \neq 1$
- Define a map of matrads $\alpha_A : C_*(KK) \rightarrow \mathcal{E}nd_{TA}$ sending $\theta_m^n \mapsto \omega_A^{n,m}$
- Define a map of A_∞ -bialgebras $G = \{g_m^n \mid g_1^1 = g\} : A \Rightarrow B$

The Transfer Algorithm: Initialization

- Define $\beta : C_0(JJ_{1,1}) \rightarrow \text{Hom}(A, B)$ by

$$f_1^1 = \text{red vertical bar} \mapsto g$$

The Transfer Algorithm: Initialization

- Define $\beta : C_0(JJ_{1,1}) \rightarrow \text{Hom}(A, B)$ by

$$f_1^1 = \text{red vertical line} \mapsto g$$

- Define β on the vertex  of $JJ_{1,2}$ by

$$\theta_2^1(f_1^1 \otimes f_1^1) = \text{Y-vertex with red legs} \mapsto \omega_B^{1,2}(g \otimes g)$$

The Transfer Algorithm: Initialization

- Define $\beta : C_0(JJ_{1,1}) \rightarrow \text{Hom}(A, B)$ by

$$f_1^1 = \text{red vertical line} \mapsto g$$

- Define β on the vertex  of $JJ_{1,2}$ by

$$\theta_2^1(f_1^1 \otimes f_1^1) = \text{red lines at bottom, black line at top} \mapsto \omega_B^{1,2}(g \otimes g)$$

- Define β on the vertex  of $JJ_{2,1}$ by

$$\theta_1^2 f_1^1 = \text{red line at bottom, black lines at top} \mapsto \omega_B^{2,1} g$$

The Transfer Algorithm: Initialization

- Consider the ∇ -cocycle $\omega_B^{1,2}(g \otimes g)$

The Transfer Algorithm: Initialization

- Consider the ∇ -cocycle $\omega_B^{1,2}(g \otimes g)$
- Choose a cocycle $\omega_A^{1,2} \in \mathcal{E}nd_{TA}$ such that

$$\tilde{g}_*[\omega_A^{1,2}] = [\omega_B^{1,2}(g \otimes g)]$$

The Transfer Algorithm: Initialization

- Consider the ∇ -cocycle $\omega_B^{1,2}(g \otimes g)$
- Choose a cocycle $\omega_A^{1,2} \in \mathcal{E}nd_{TA}$ such that

$$\tilde{g}_*[\omega_A^{1,2}] = [\omega_B^{1,2}(g \otimes g)]$$

- Define $\alpha_A : C_0(KK_{1,2}) \rightarrow Hom(A^{\otimes 2}, A)$ by $\bigwedge \mapsto \omega_A^{1,2}$

The Transfer Algorithm: Initialization

- Consider the ∇ -cocycle $\omega_B^{1,2}(g \otimes g)$
- Choose a cocycle $\omega_A^{1,2} \in \mathcal{E}nd_{TA}$ such that

$$\tilde{g}_*[\omega_A^{1,2}] = [\omega_B^{1,2}(g \otimes g)]$$

- Define $\alpha_A : C_0(KK_{1,2}) \rightarrow \text{Hom}(A^{\otimes 2}, A)$ by $\text{Y} \mapsto \omega_A^{1,2}$
- Define $\alpha_A : C_0(\partial KK_{1,3}) \rightarrow \text{Hom}(A^{\otimes 3}, A)$ by

$$\text{Y} \mapsto \omega_A^{1,2}(\omega_A^{1,2} \otimes \mathbf{1}) \text{ and } \text{Y} \mapsto \omega_A^{1,2}(\mathbf{1} \otimes \omega_A^{1,2})$$

The Transfer Algorithm: Initialization

- Consider the ∇ -cocycle $\omega_B^{1,2}(g \otimes g)$
- Choose a cocycle $\omega_A^{1,2} \in \mathcal{E}nd_{TA}$ such that

$$\tilde{g}_*[\omega_A^{1,2}] = [\omega_B^{1,2}(g \otimes g)]$$

- Define $\alpha_A : C_0(KK_{1,2}) \rightarrow \text{Hom}(A^{\otimes 2}, A)$ by $\text{Y} \mapsto \omega_A^{1,2}$
- Define $\alpha_A : C_0(\partial KK_{1,3}) \rightarrow \text{Hom}(A^{\otimes 3}, A)$ by

$$\text{Y} \mapsto \omega_A^{1,2}(\omega_A^{1,2} \otimes \mathbf{1}) \text{ and } \text{Y} \mapsto \omega_A^{1,2}(\mathbf{1} \otimes \omega_A^{1,2})$$

- Extend β to the vertex $\text{Y} \subset JJ_{1,2}$ via $\text{Y} \mapsto g\omega_A^{1,2}$

The Transfer Algorithm: Initialization

- Dually, consider the ∇ -cocycle $\omega_B^{2,1}g$

The Transfer Algorithm: Initialization

- Dually, consider the ∇ -cocycle $\omega_B^{2,1}g$
- Choose a cocycle $\omega_A^{2,1} \in \mathcal{E}nd_{TA}$ such that

$$\tilde{g}_*[\omega_A^{2,1}] = [\omega_B^{2,1}g]$$

The Transfer Algorithm: Initialization

- Dually, consider the ∇ -cocycle $\omega_B^{2,1}g$
- Choose a cocycle $\omega_A^{2,1} \in \mathcal{E}nd_{TA}$ such that

$$\tilde{g}_*[\omega_A^{2,1}] = [\omega_B^{2,1}g]$$

- Define $\alpha_A : C_0(KK_{2,1}) \rightarrow Hom(A, A^{\otimes 2})$ by $\bigvee \mapsto \omega_A^{2,1}$

The Transfer Algorithm: Initialization

- Dually, consider the ∇ -cocycle $\omega_B^{2,1}g$
- Choose a cocycle $\omega_A^{2,1} \in \mathcal{E}nd_{TA}$ such that

$$\tilde{g}_*[\omega_A^{2,1}] = [\omega_B^{2,1}g]$$

- Define $\alpha_A : C_0(KK_{2,1}) \rightarrow Hom(A, A^{\otimes 2})$ by $\Upsilon \mapsto \omega_A^{2,1}$
- Define $\alpha_A : C_0(\partial KK_{3,1}) \rightarrow Hom(A, A^{\otimes 3})$ by

$$\Upsilon \mapsto (\omega_A^{2,1} \otimes \mathbf{1}) \omega_A^{2,1} \text{ and } \Upsilon \mapsto (\mathbf{1} \otimes \omega_A^{2,1}) \omega_A^{2,1}$$

The Transfer Algorithm: Initialization

- Dually, consider the ∇ -cocycle $\omega_B^{2,1}g$
- Choose a cocycle $\omega_A^{2,1} \in \mathcal{E}nd_{TA}$ such that

$$\tilde{g}_*[\omega_A^{2,1}] = [\omega_B^{2,1}g]$$

- Define $\alpha_A : C_0(KK_{2,1}) \rightarrow Hom(A, A^{\otimes 2})$ by $\text{Y} \mapsto \omega_A^{2,1}$
- Define $\alpha_A : C_0(\partial KK_{3,1}) \rightarrow Hom(A, A^{\otimes 3})$ by

$$\text{Y} \mapsto (\omega_A^{2,1} \otimes \mathbf{1}) \omega_A^{2,1} \text{ and } \text{Y} \mapsto (\mathbf{1} \otimes \omega_A^{2,1}) \omega_A^{2,1}$$

- Extend β to the vertex $\text{Y} \subset JJ_{1,2}$ via $\text{Y} \mapsto (g \otimes g) \omega_A^{2,1}$

The Transfer Algorithm: Initialization

- Define $\alpha_A : C_0(\partial KK_{2,2}) \rightarrow \text{Hom}(A^{\otimes 2}, A^{\otimes 2})$ by

$$\text{Diagram} \mapsto \left(\omega_A^{1,2} \otimes \omega_A^{1,2} \right) \sigma_{2,2} \left(\omega_A^{2,1} \otimes \omega_A^{2,1} \right)$$

and

$$\text{Diagram} \mapsto \omega_A^{2,1} \omega_A^{1,2}$$

The Transfer Algorithm: Initialization

- Note that $\left[\omega_B^{1,2} (g \otimes g) - g \omega_A^{1,2} \right] = 0$

The Transfer Algorithm: Initialization

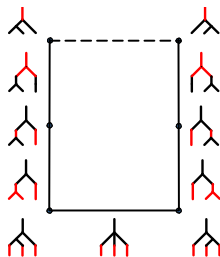
- Note that $\left[\omega_B^{1,2} (g \otimes g) - g \omega_A^{1,2} \right] = 0$
- Choose a cochain g_2^1 such that $\nabla g_2^1 = \omega_B^{1,2} (g \otimes g) - g \omega_A^{1,2}$

The Transfer Algorithm: Initialization

- Note that $\left[\omega_B^{1,2} (g \otimes g) - g \omega_A^{1,2} \right] = 0$
- Choose a cochain g_2^1 such that $\nabla g_2^1 = \omega_B^{1,2} (g \otimes g) - g \omega_A^{1,2}$
- Define $\beta : C_1(JJ_{1,2}) \rightarrow \text{Hom}(A^{\otimes 2}, B)$ by $f_2^1 = \text{red trivalent vertex} \mapsto g_2^1$

The Transfer Algorithm: Initialization

- Note that $\left[\omega_B^{1,2} (g \otimes g) - g \omega_A^{1,2} \right] = 0$
- Choose a cochain g_2^1 such that $\nabla g_2^1 = \omega_B^{1,2} (g \otimes g) - g \omega_A^{1,2}$
- Define $\beta : C_1(JJ_{1,2}) \rightarrow \text{Hom}(A^{\otimes 2}, B)$ by $f_2^1 = \text{red trivalent vertex} \mapsto g_2^1$
- Define β on $\partial JJ_{1,3} \setminus \text{int } KK_{1,3} \times 1$:



The Transfer Algorithm: Initialization

- Dually, note that $\left[\omega_B^{2,1} g - (g \otimes g) \omega_A^{2,1} \right] = 0$

The Transfer Algorithm: Initialization

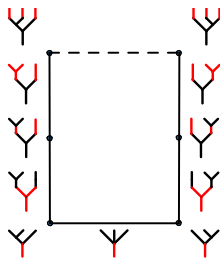
- Dually, note that $\left[\omega_B^{2,1} g - (g \otimes g) \omega_A^{2,1} \right] = 0$
- Choose a cochain $g_1^2 \in U_{A,B}$ such that $\nabla g_1^2 = \omega_B^{2,1} g - (g \otimes g) \omega_A^{2,1}$

The Transfer Algorithm: Initialization

- Dually, note that $\left[\omega_B^{2,1} g - (g \otimes g) \omega_A^{2,1} \right] = 0$
- Choose a cochain $g_1^2 \in U_{A,B}$ such that $\nabla g_1^2 = \omega_B^{2,1} g - (g \otimes g) \omega_A^{2,1}$
- Define $\beta : C_1(JJ_{2,1}) \rightarrow \text{Hom}(A, B^{\otimes 2})$ by $f_1^2 = \textcolor{red}{Y} \mapsto g_1^2$

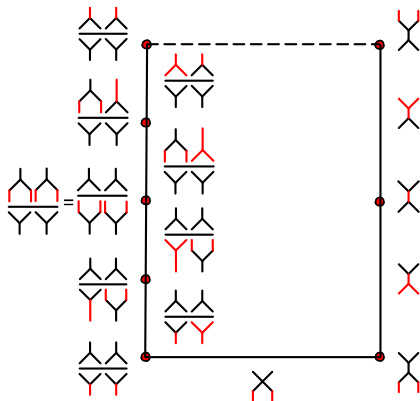
The Transfer Algorithm: Initialization

- Dually, note that $\left[\omega_B^{2,1} g - (g \otimes g) \omega_A^{2,1} \right] = 0$
- Choose a cochain $g_1^2 \in U_{A,B}$ such that $\nabla g_1^2 = \omega_B^{2,1} g - (g \otimes g) \omega_A^{2,1}$
- Define $\beta : C_1(JJ_{2,1}) \rightarrow \text{Hom}(A, B^{\otimes 2})$ by $f_1^2 = \text{Y} \mapsto g_1^2$
- Define β on $\partial JJ_{3,1} \setminus \text{int } KK_{3,1} \times 1$:



The Transfer Algorithm: Initialization

Having defined β on $\text{red trivalent vertex}$ and red Y-vertex , define β on $\partial JJ_{2,2} \setminus \text{int } KK_{2,2} \times 1$:



The Transfer Algorithm: Induction Hypothesis

Given $m + n \geq 4$, assume for $i + j < m + n$, $ij \neq 1$, there exist maps

- $\alpha_A : C_*(KK_{j,i}) \rightarrow \text{Hom}(A^{\otimes i}, A^{\otimes j})$ of matrads sending $\theta_i^j \mapsto \omega_A^{j,i}$

The Transfer Algorithm: Induction Hypothesis

Given $m + n \geq 4$, assume for $i + j < m + n$, $ij \neq 1$, there exist maps

- $\alpha_A : C_*(KK_{j,i}) \rightarrow \text{Hom}(A^{\otimes i}, A^{\otimes j})$ of matrads sending $\theta_i^j \mapsto \omega_A^{j,i}$
- $\beta : C_*(JJ_{j,i}) \rightarrow \text{Hom}(A^{\otimes i}, B^{\otimes j})$ of rel. matrads sending $f_i^j \mapsto g_i^j$

The Transfer Algorithm: Induction Objectives

- Define α_A on the generator $\theta_m^n \in C_{m+n-3}(KK_{n,m})$

The Transfer Algorithm: Induction Objectives

- Define α_A on the generator $\theta_m^n \in C_{m+n-3}(KK_{n,m})$
- Define β on the monomial $(f_1^1)^{\otimes n} \theta_m^n \in C_{m+n-3}(JJ_{n,m})$

The Transfer Algorithm: Induction Objectives

- Define α_A on the generator $\theta_m^n \in C_{m+n-3}(KK_{n,m})$
- Define β on the monomial $(f_1^1)^{\otimes n} \theta_m^n \in C_{m+n-3}(JJ_{n,m})$
- Define β on the generator $f_m^n \in C_{m+n-2}(JJ_{n,m})$

The Transfer Algorithm: Induction

For each $i + j = m + n$, $ij \neq 1$

- Define α_A on a monomial in $C_*(\partial KK_{n,m})$ to be its corresponding fraction product of operations in $\{\omega_A^{j,i}\}$; let

$$z = \alpha_A (C_{m+n-2}(\partial KK_{n,m}))$$

The Transfer Algorithm: Induction

For each $i + j = m + n$, $ij \neq 1$

- Define α_A on a monomial in $C_*(\partial KK_{n,m})$ to be its corresponding fraction product of operations in $\{\omega_A^{j,i}\}$; let

$$z = \alpha_A (C_{m+n-2}(\partial KK_{n,m}))$$

- Define β on a monomial in $C_*(\partial JJ_{n,m} \setminus \text{int } KK_{n,m} \times 1)$ to be its corresponding fraction product of operations/maps in $\{\omega_A^{j,i}, g_i^j, \omega_B^{j,i}\}$; let

$$\varphi = \beta (C_{m+n-1}(\partial JJ_{n,m} \setminus \text{int } KK_{n,m} \times 1))$$

The Transfer Algorithm: Induction

For each $i + j = m + n$, $ij \neq 1$

- Define α_A on a monomial in $C_*(\partial KK_{n,m})$ to be its corresponding fraction product of operations in $\{\omega_A^{j,i}\}$; let

$$z = \alpha_A (C_{m+n-2}(\partial KK_{n,m}))$$

- Define β on a monomial in $C_*(\partial JJ_{n,m} \setminus \text{int } KK_{n,m} \times 1)$ to be its corresponding fraction product of operations/maps in $\{\omega_A^{j,i}, g_i^j, \omega_B^{j,i}\}$; let

$$\varphi = \beta (C_{m+n-1}(\partial JJ_{n,m} \setminus \text{int } KK_{n,m} \times 1))$$

- Then $\tilde{g}(z) = \nabla \varphi$ implies $[z] = 0$

The Transfer Algorithm: Induction

For each $i + j = m + n$, $ij \neq 1$

- Define α_A on a monomial in $C_*(\partial KK_{n,m})$ to be its corresponding fraction product of operations in $\{\omega_A^{j,i}\}$; let

$$z = \alpha_A (C_{m+n-2}(\partial KK_{n,m}))$$

- Define β on a monomial in $C_*(\partial JJ_{n,m} \setminus \text{int } KK_{n,m} \times 1)$ to be its corresponding fraction product of operations/maps in $\{\omega_A^{j,i}, g_i^j, \omega_B^{j,i}\}$; let

$$\varphi = \beta (C_{m+n-1}(\partial JJ_{n,m} \setminus \text{int } KK_{n,m} \times 1))$$

- Then $\tilde{g}(z) = \nabla \varphi$ implies $[z] = 0$
- Choose a cochain b such that $\nabla b = z$

The Transfer Algorithm: Induction

- Note that $\nabla (\tilde{g}(b) - \varphi) = \nabla \tilde{g}(b) - \tilde{g}(z) = \tilde{g}(\nabla b - z) = 0$

The Transfer Algorithm: Induction

- Note that $\nabla (\tilde{g}(b) - \varphi) = \nabla \tilde{g}(b) - \tilde{g}(z) = \tilde{g}(\nabla b - z) = 0$
- Choose a cocycle $u \in \tilde{g}_*^{-1} [\tilde{g}(b) - \varphi]$

The Transfer Algorithm: Induction

- Note that $\nabla (\tilde{g}(b) - \varphi) = \nabla \tilde{g}(b) - \tilde{g}(z) = \tilde{g}(\nabla b - z) = 0$
- Choose a cocycle $u \in \tilde{g}_*^{-1} [\tilde{g}(b) - \varphi]$
- Define $\alpha_A(\theta_m^n) = \omega_A^{n,m} := b - u$

The Transfer Algorithm: Induction

- Note that $\nabla (\tilde{g}(b) - \varphi) = \nabla \tilde{g}(b) - \tilde{g}(z) = \tilde{g}(\nabla b - z) = 0$
- Choose a cocycle $u \in \tilde{g}_*^{-1} [\tilde{g}(b) - \varphi]$
- Define $\alpha_A(\theta_m^n) = \omega_A^{n,m} := b - u$
- Define $\beta \left((f_1^1)^{\otimes n} \theta_m^n \right) = g^{\otimes n} \omega_A^{n,m}$

The Transfer Algorithm: The Induction

- Note that $[\tilde{g}(\omega_A^{n,m}) - \varphi] = [\tilde{g}(b - u) - \varphi]$
$$= [\tilde{g}(b) - \varphi] - [\tilde{g}(u)] = 0$$

The Transfer Algorithm: The Induction

- Note that $[\tilde{g}(\omega_A^{n,m}) - \varphi] = [\tilde{g}(b - u) - \varphi]$
$$= [\tilde{g}(b) - \varphi] - [\tilde{g}(u)] = 0$$

- Choose a cochain g_m^n such that

$$\nabla g_m^n = \tilde{g}(\omega_A^{n,m}) - \varphi = g^{\otimes n} \omega_A^{n,m} - \varphi$$

The Transfer Algorithm: The Induction

- Note that $[\tilde{g}(\omega_A^{n,m}) - \varphi] = [\tilde{g}(b - u) - \varphi]$

$$= [\tilde{g}(b) - \varphi] - [\tilde{g}(u)] = 0$$

- Choose a cochain g_m^n such that

$$\nabla g_m^n = \tilde{g}(\omega_A^{n,m}) - \varphi = g^{\otimes n} \omega_A^{n,m} - \varphi$$

- Define $\beta(f_m^n) = g_m^n$

The Transfer Algorithm: The Induction

- Note that $[\tilde{g}(\omega_A^{n,m}) - \varphi] = [\tilde{g}(b - u) - \varphi]$
$$= [\tilde{g}(b) - \varphi] - [\tilde{g}(u)] = 0$$

- Choose a cochain g_m^n such that

$$\nabla g_m^n = \tilde{g}(\omega_A^{n,m}) - \varphi = g^{\otimes n} \omega_A^{n,m} - \varphi$$

- Define $\beta(f_m^n) = g_m^n$
- This completes the induction

Application: Homology of Loop Spaces

- Let $\mathbf{k}$ be a field and let X be a space

Application: Homology of Loop Spaces

- Let $\mathbf{k}$ be a field and let X be a space
- Consider the Hopf algebra $S_*(\Omega X; \mathbf{k})$ of singular chains on Moore (base-pointed) loops

Application: Homology of Loop Spaces

- Let $\mathbf{k}$ be a field and let X be a space
- Consider the Hopf algebra $S_*(\Omega X; \mathbf{k})$ of singular chains on Moore (base-pointed) loops
- **Theorem.** $H_*(\Omega X; \mathbf{k})$ is an A_∞ -bialgebra

Application: Homology of Loop Spaces

- Let $\mathbf{k}$ be a field and let X be a space
- Consider the Hopf algebra $S_*(\Omega X; \mathbf{k})$ of singular chains on Moore (base-pointed) loops
- **Theorem.** $H_*(\Omega X; \mathbf{k})$ is an A_∞ -bialgebra
- *Proof.* A cycle-selecting chain map $g : H_*(\Omega X; \mathbf{k}) \rightarrow S_*(\Omega X; \mathbf{k})$ induces a homology isomorphism $\tilde{g}$ by Prop 1

Application: Homology of Loop Spaces

- Let $\mathbf{k}$ be a field and let X be a space
- Consider the Hopf algebra $S_*(\Omega X; \mathbf{k})$ of singular chains on Moore (base-pointed) loops
- **Theorem.** $H_*(\Omega X; \mathbf{k})$ is an A_∞ -bialgebra
- *Proof.* A cycle-selecting chain map $g : H_*(\Omega X; \mathbf{k}) \rightarrow S_*(\Omega X; \mathbf{k})$ induces a homology isomorphism $\tilde{g}$ by Prop 1
- The Transfer Thm imposes an A_∞ -bialgebra structure on $H(\Omega X; \mathbf{k})$

Example

- $Y = (S^2 \times S^3) \vee \Sigma \mathbb{C}P^2$

Example

- $Y = (S^2 \times S^3) \vee \Sigma \mathbb{C}P^2$
- Mult generators $\bar{a}_i \in H^i(S^i; \mathbb{Z}_2)$ and $\bar{b} \in H^3(\Sigma \mathbb{C}P^2; \mathbb{Z}_2)$

Example

- $Y = (S^2 \times S^3) \vee \Sigma \mathbb{C}P^2$
- Mult generators $\bar{a}_i \in H^i(S^i; \mathbb{Z}_2)$ and $\bar{b} \in H^3(\Sigma \mathbb{C}P^2; \mathbb{Z}_2)$
- Let $f : Y \rightarrow K(\mathbb{Z}_2, 5)$ be a map such that $f^*(\iota_5) = \bar{a}_2 \bar{a}_3 + Sq^2 \bar{b}$

Example

- $Y = (S^2 \times S^3) \vee \Sigma \mathbb{C}P^2$
- Mult generators $\bar{a}_i \in H^i(S^i; \mathbb{Z}_2)$ and $\bar{b} \in H^3(\Sigma \mathbb{C}P^2; \mathbb{Z}_2)$
- Let $f : Y \rightarrow K(\mathbb{Z}_2, 5)$ be a map such that $f^*(\iota_5) = \bar{a}_2 \bar{a}_3 + Sq^2 \bar{b}$
- Consider the fibration $p : X \rightarrow Y$ induced by f

Example

- $Y = (S^2 \times S^3) \vee \Sigma \mathbb{C}P^2$
- Mult generators $\bar{a}_i \in H^i(S^i; \mathbb{Z}_2)$ and $\bar{b} \in H^3(\Sigma \mathbb{C}P^2; \mathbb{Z}_2)$
- Let $f : Y \rightarrow K(\mathbb{Z}_2, 5)$ be a map such that $f^*(\iota_5) = \bar{a}_2 \bar{a}_3 + Sq^2 \bar{b}$
- Consider the fibration $p : X \rightarrow Y$ induced by f
- $\bar{a}_i \mapsto a_i$ and $\bar{b} \mapsto b$ under $p^* : H^*(Y; \mathbb{Z}_2) \rightarrow H^*(X; \mathbb{Z}_2)$

Example

- $Y = (S^2 \times S^3) \vee \Sigma \mathbb{C}P^2$
- Mult generators $\bar{a}_i \in H^i(S^i; \mathbb{Z}_2)$ and $\bar{b} \in H^3(\Sigma \mathbb{C}P^2; \mathbb{Z}_2)$
- Let $f : Y \rightarrow K(\mathbb{Z}_2, 5)$ be a map such that $f^*(\iota_5) = \bar{a}_2 \bar{a}_3 + Sq^2 \bar{b}$
- Consider the fibration $p : X \rightarrow Y$ induced by f
- $\bar{a}_i \mapsto a_i$ and $\bar{b} \mapsto b$ under $p^* : H^*(Y; \mathbb{Z}_2) \rightarrow H^*(X; \mathbb{Z}_2)$
- $A = H^*(X; \mathbb{Z}_2) = \{1, a_2, a_3, b, a_2 a_3 = Sq^2 b, \dots\}$

A homotopy Gerstenhaber algebra structure on A

- Induced by the attaching map in $\mathbb{C}P^2$ and the twisting in X :

$$E_{p,q} : A^{\otimes p} \otimes A^{\otimes q} \rightarrow A$$

A homotopy Gerstenhaber algebra structure on A

- Induced by the attaching map in $\mathbb{C}P^2$ and the twisting in X :

$$E_{p,q} : A^{\otimes p} \otimes A^{\otimes q} \rightarrow A$$

- Nontrivial action $E_{1,0} = E_{0,1} = \text{Id}_A$ and $E_{1,1}(b; b) = a_2 a_3$

A homotopy Gerstenhaber algebra structure on A

- Induced by the attaching map in $\mathbb{C}P^2$ and the twisting in X :

$$E_{p,q} : A^{\otimes p} \otimes A^{\otimes q} \rightarrow A$$

- Nontrivial action $E_{1,0} = E_{0,1} = \text{Id}_A$ and $E_{1,1}(b; b) = a_2 a_3$
- Consider the Bar Construction BA with cofree coproduct Δ

A homotopy Gerstenhaber algebra structure on A

- Induced by the attaching map in $\mathbb{C}P^2$ and the twisting in X :

$$E_{p,q} : A^{\otimes p} \otimes A^{\otimes q} \rightarrow A$$

- Nontrivial action $E_{1,0} = E_{0,1} = \text{Id}_A$ and $E_{1,1}(b; b) = a_2 a_3$
- Consider the Bar Construction BA with cofree coproduct Δ
- And the tensor product coalgebra $BA \otimes BA$ with $\psi = \sigma_{2,2}(\Delta \otimes \Delta)$

A homotopy Gerstenhaber algebra structure on A

- Induced by the attaching map in $\mathbb{C}P^2$ and the twisting in X :

$$E_{p,q} : A^{\otimes p} \otimes A^{\otimes q} \rightarrow A$$

- Nontrivial action $E_{1,0} = E_{0,1} = \text{Id}_A$ and $E_{1,1}(b; b) = a_2 a_3$
- Consider the Bar Construction BA with cofree coproduct Δ
- And the tensor product coalgebra $BA \otimes BA$ with $\psi = \sigma_{2,2}(\Delta \otimes \Delta)$
- Induced map $\phi = E'_{1,0} + E'_{0,1} + E'_{1,1} : BA \otimes BA \rightarrow A$ acts nontrivially:

$$\phi([b] \otimes [b]) = a_2 a_3$$

Induced multiplication on BA

- ϕ lifts to a chain map $\mu : BA \otimes BA \rightarrow BA$ of DG coalgebras given by

$$\mu = \sum_{k \geq 0} \downarrow^{\otimes k+1} \phi^{\otimes k+1} \bar{\psi}^{(k)}$$

where $\bar{\psi}^{(0)} = \text{Id}_{BA \otimes BA}$ and $\bar{\psi}^{(k)}$ is the left-iterated reduced coproduct

Induced multiplication on BA

- ϕ lifts to a chain map $\mu : BA \otimes BA \rightarrow BA$ of DG coalgebras given by

$$\mu = \sum_{k \geq 0} \downarrow^{\otimes k+1} \phi^{\otimes k+1} \bar{\psi}^{(k)}$$

where $\bar{\psi}^{(0)} = \text{Id}_{BA \otimes BA}$ and $\bar{\psi}^{(k)}$ is the left-iterated reduced coproduct

- $(BA, d_{BA}, \Delta, \mu)$ is a DG Hopf algebra

Induced multiplication on BA

- ϕ lifts to a chain map $\mu : BA \otimes BA \rightarrow BA$ of DG coalgebras given by

$$\mu = \sum_{k \geq 0} \downarrow^{\otimes k+1} \phi^{\otimes k+1} \bar{\psi}^{(k)}$$

where $\bar{\psi}^{(0)} = \text{Id}_{BA \otimes BA}$ and $\bar{\psi}^{(k)}$ is the left-iterated reduced coproduct

- $(BA, d_{BA}, \Delta, \mu)$ is a DG Hopf algebra
- Nontrivial product $\mu([b] \otimes [b]) = [a_2 a_3]$

Induced DG Hopf algebra structure on BA

- $H = H^*(BA) \approx H^*(\Omega X; \mathbb{Z}_2)$ as graded coalgebras

Induced DG Hopf algebra structure on BA

- $H = H^*(BA) \approx H^*(\Omega X; \mathbb{Z}_2)$ as graded coalgebras
- Denote: $\alpha_i = \text{cls}[a_i] \quad \beta = \text{cls}[b]$

Induced DG Hopf algebra structure on BA

- $H = H^*(BA) \approx H^*(\Omega X; \mathbb{Z}_2)$ as graded coalgebras
- Denote: $\alpha_i = \text{cls}[a_i] \quad \beta = \text{cls}[b]$
- Choose a cycle-selecting map $g : H \rightarrow A$ such that

$$g(\text{cls}[x_1 | \cdots | x_n]) = [x_1 | \cdots | x_n]$$

g is homotopy multiplicative

- Let $\varphi : H \otimes H \rightarrow H$ be the multiplication induced by μ

g is homotopy multiplicative

- Let $\varphi : H \otimes H \rightarrow H$ be the multiplication induced by μ
- $\mu([b] \otimes [b]) = [a_2 a_3] = d_{BA}[a_2 | a_3]$ implies $\varphi(\beta \otimes \beta) = 0$

g is homotopy multiplicative

- Let $\varphi : H \otimes H \rightarrow H$ be the multiplication induced by μ
- $\mu([b] \otimes [b]) = [a_2 a_3] = d_{BA}[a_2 | a_3]$ implies $\varphi(\beta \otimes \beta) = 0$
- The Transfer Theorem implies the relation on $JJ_{1,2}$ holds:

$$\nabla g_2 = g\varphi + \mu(g \otimes g)$$

g is homotopy multiplicative

- Let $\varphi : H \otimes H \rightarrow H$ be the multiplication induced by μ
- $\mu([b] \otimes [b]) = [a_2 a_3] = d_{BA} [a_2 | a_3]$ implies $\varphi(\beta \otimes \beta) = 0$
- The Transfer Theorem implies the relation on $JJ_{1,2}$ holds:

$$\nabla g_2 = g\varphi + \mu(g \otimes g)$$

- $(g\varphi + \mu(g \otimes g))(\beta \otimes \beta) = [a_2 a_3] = d_{BA} [a_2 | a_3]$ implies

g is homotopy multiplicative

- Let $\varphi : H \otimes H \rightarrow H$ be the multiplication induced by μ
- $\mu([b] \otimes [b]) = [a_2 a_3] = d_{BA} [a_2 | a_3]$ implies $\varphi(\beta \otimes \beta) = 0$
- The Transfer Theorem implies the relation on $JJ_{1,2}$ holds:

$$\nabla g_2 = g\varphi + \mu(g \otimes g)$$

- $(g\varphi + \mu(g \otimes g))(\beta \otimes \beta) = [a_2 a_3] = d_{BA} [a_2 | a_3]$ implies
- *There is a homotopy g_2 satisfying the relation above such that*

$$g_2(\beta \otimes \beta) = [a_2 | a_3]$$

g is homotopy comultiplicative

- Let $\Delta_H : H \rightarrow H \otimes H$ be the induced coproduct

g is homotopy comultiplicative

- Let $\Delta_H : H \rightarrow H \otimes H$ be the induced coproduct
- The Transfer Theorem implies the relation on $JJ_{2,1}$ holds:

$$\nabla g^2 = \Delta g + (g \otimes g) \Delta_H$$

g is homotopy comultiplicative

- Let $\Delta_H : H \rightarrow H \otimes H$ be the induced coproduct
- The Transfer Theorem implies the relation on $JJ_{2,1}$ holds:

$$\nabla g^2 = \Delta g + (g \otimes g) \Delta_H$$

- β primitive implies

$$[\Delta g + (g \otimes g) \Delta_H] (\beta) = 0$$

g is homotopy comultiplicative

- Let $\Delta_H : H \rightarrow H \otimes H$ be the induced coproduct
- The Transfer Theorem implies the relation on $JJ_{2,1}$ holds:

$$\nabla g^2 = \Delta g + (g \otimes g) \Delta_H$$

- β primitive implies

$$[\Delta g + (g \otimes g) \Delta_H] (\beta) = 0$$

- *There is a homotopy g^2 satisfying the relation on $JJ_{2,1}$ such that*

$$g^2 (\beta) = 0$$

Non-operadic operation defined

- The Transfer Theorem implies the relation on $JJ_{2,2}$ holds:

$$\begin{aligned}\nabla g_2^2 &= (\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi \\ &\quad + (\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2 \\ &\quad + (g \otimes g) \omega_2^2\end{aligned}$$

Non-operadic operation defined

- The Transfer Theorem implies the relation on $JJ_{2,2}$ holds:

$$\begin{aligned}\nabla g_2^2 &= (\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi \\ &\quad + (\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2 \\ &\quad + (g \otimes g) \omega_2^2\end{aligned}$$

- $((\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi) (\beta \otimes \beta) = 0$

Non-operadic operation defined

- The Transfer Theorem implies the relation on $JJ_{2,2}$ holds:

$$\begin{aligned}\nabla g_2^2 &= (\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi \\ &\quad + (\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2 \\ &\quad + (g \otimes g) \omega_2^2\end{aligned}$$

- $((\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi) (\beta \otimes \beta) = 0$
- $((\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2) (\beta \otimes \beta)$
 $= [a_2] \otimes [a_3]$

Non-operadic operation defined

- The Transfer Theorem implies the relation on $JJ_{2,2}$ holds:

$$\begin{aligned}\nabla g_2^2 &= (\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi \\ &\quad + (\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2 \\ &\quad + (g \otimes g) \omega_2^2\end{aligned}$$

- $((\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi) (\beta \otimes \beta) = 0$
- $((\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2) (\beta \otimes \beta)$
 $= [a_2] \otimes [a_3]$
- $[a_2] \otimes [a_3]$ represents a non-zero class in $BA \otimes BA$

Non-operadic operation defined

- The Transfer Theorem implies the relation on $JJ_{2,2}$ holds:

$$\begin{aligned}\nabla g_2^2 &= (\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi \\ &\quad + (\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2 \\ &\quad + (g \otimes g) \omega_2^2\end{aligned}$$

- $((\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi) (\beta \otimes \beta) = 0$
- $((\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2) (\beta \otimes \beta)$
 $= [a_2] \otimes [a_3]$
- $[a_2] \otimes [a_3]$ represents a non-zero class in $BA \otimes BA$
- Thus no choice of g_2^2 kills the obstruction $[a_2] \otimes [a_3]$

Conclusion:

- *There is an operation ω_2^2 and a cochain homotopy g_2^2 satisfying the relation on $JJ_{2,2}$ such that*

$$\omega_2^2(\beta \otimes \beta) = \alpha_2 \otimes \alpha_3$$

Conclusion:

- *There is an operation ω_2^2 and a cochain homotopy g_2^2 satisfying the relation on $JJ_{2,2}$ such that*

$$\omega_2^2(\beta \otimes \beta) = \alpha_2 \otimes \alpha_3$$

- Then $(g \otimes g) \omega_2^2(\beta \otimes \beta) = [a_2] \otimes [a_3]$ kills the obstruction

Conclusion:

- *There is an operation ω_2^2 and a cochain homotopy g_2^2 satisfying the relation on $JJ_{2,2}$ such that*

$$\omega_2^2(\beta \otimes \beta) = \alpha_2 \otimes \alpha_3$$

- Then $(g \otimes g) \omega_2^2(\beta \otimes \beta) = [a_2] \otimes [a_3]$ kills the obstruction
- We have identified the first non-operadic cohomology operation

Thank you!