

Matradic Operations on Loop Cohomology

Joint work with Samson Saneblidze

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Millersville Univ of Pennsylvania

U. Penn Deformation Theory Seminar

August 18, 2011

Goal of the talk

- To identify a space X such that $H = H^*(\Omega X; \mathbb{Z}_2)$ comes equipped with an induced operation

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- We assume existence of biassociahedra $KK = \{KK_{n,m}\}$ and bmultipliheda $JJ = \{JJ_{n,m}\}$
- The operation ω_2^2 will be defined in terms of $JJ_{2,1} = JJ_{1,2}$ and $JJ_{2,2}$

The Free Matrad

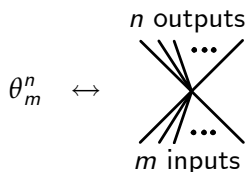
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- $\mathcal{H}_\infty$ is realized by cellular chains $C_*(KK)$
- Top dim'l cell of $KK_{n,m}$ is identified with the matradic generator



The Free Matrad

- Cells of KK are identified with certain fraction product monomials

The diagram shows an equation between three matradic diagrams. On the left, a fraction where the numerator consists of two separate vertices (each with one incoming line from the top and two outgoing lines to the bottom) and the denominator consists of three separate vertices (each with two incoming lines from the top and one outgoing line to the bottom). This is followed by an equals sign and a single matradic diagram. This diagram is a diamond shape with a vertical line through the center, and a diagonal line from the top-left to the bottom-right. It has four external lines: one at the top, one at the bottom, one on the left, and one on the right. To the right of this diagram is a double-headed arrow pointing to the text "a vertex of $KK_{2,3}$ ".

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$$\frac{\begin{array}{cc} \text{Y-shape} & \text{Y-shape} \\ \hline \text{Y-shape} & \text{Y-shape} & \text{Y-shape} \end{array}}{=} \text{A complex diagram with multiple Y-shapes and lines} \leftrightarrow \text{a vertex of } KK_{2,3}$$

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- $KK_{n,1} = KK_{1,n} = K_n$ (Stasheff's associahedron)
- $\dim(KK_{n,m}) = m + n - 3$

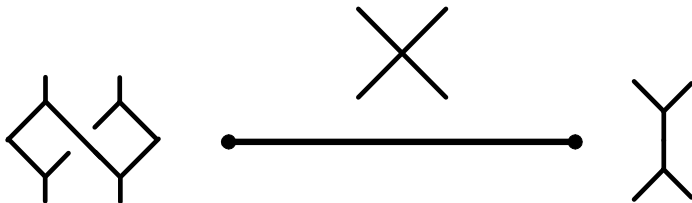
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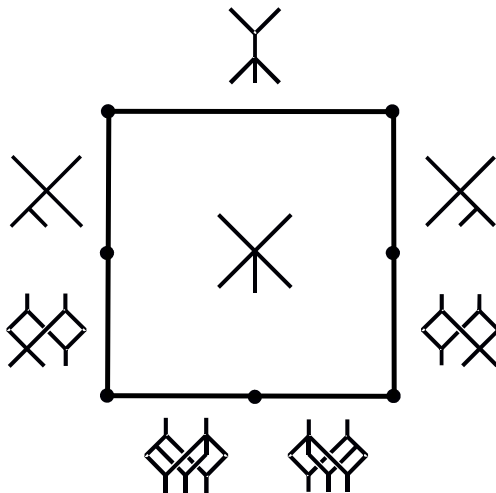
$$\frac{\begin{array}{cc} \text{Y-shape} & \text{Y-shape} \\ \hline \text{Y-shape} & \text{Y-shape} & \text{Y-shape} \end{array}}{=} \begin{array}{c} \text{A complex graph with 6 vertices and 9 edges, consisting of two diamond shapes sharing a central vertex and having additional internal edges.} \end{array} \leftrightarrow \text{a vertex of } KK_{2,3}$$

- $KK_{n,1} = KK_{1,n} = K_n$ (Stasheff's associahedron)
- $\dim(KK_{n,m}) = m + n - 3$
- $KK_{n,m}$ is contractible

The Biassociahedron $KK(2,2)$:



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A-infinity Bialgebras

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$$\omega = \{ \omega_m^n \in \text{Hom}(A^{\otimes m}, A^{\otimes n}) \mid mn > 1 \}$$

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- An A_∞ -bialgebra (A, d, ω) is an algebra over $\mathcal{H}_\infty$

The Free Relative Matrad

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- $\mathcal{J}\mathcal{J}_\infty$ is an $\mathcal{H}_\infty$ -bimodule generated by $\{\mathfrak{f}_m^n\}$
- $\mathcal{J}\mathcal{J}_\infty$ is realized by cellular chains $C_*(JJ)$

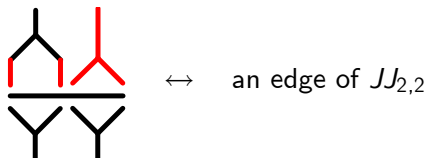
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- $\mathcal{JJ}_\infty$ is an $\mathcal{H}_\infty$ -bimodule generated by $\{\mathfrak{f}_m^n\}$
- $\mathcal{JJ}_\infty$ is realized by cellular chains $C_*(JJ)$
- Top dim'l cell of $JJ_{n,m}$ is identified with the $\mathcal{H}_\infty$ -bimodule generator

$$\mathfrak{f}_m^n \leftrightarrow \begin{array}{c} n \\ \diagup \quad \diagdown \\ \text{...} \\ \diagdown \quad \diagup \\ m \end{array}$$

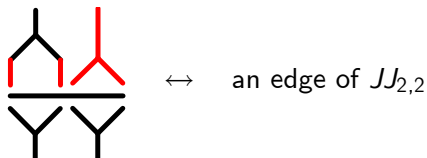
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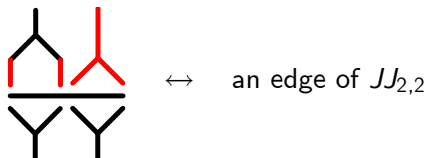
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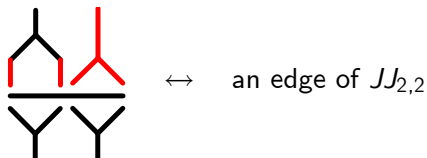
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- $JJ_{n,1} = JJ_{1,n} = J_n$ (the multiplihedron)
- $\dim(JJ_{n,m}) = m + n - 2$
- $JJ_{n,m}$ is a subdivision of $KK_{n,m} \times I$

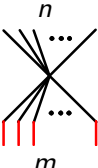
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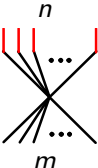
$KK_{n,m}$ embeds in $JJ_{n,m}$ as a codim 1 cell two ways:

$$\bullet \text{ As } KK_{n,m} \times 0 \iff \begin{array}{c} n \\ \diagup \quad \diagdown \\ \vdots \quad \vdots \\ \diagdown \quad \diagup \\ m \end{array} \iff \theta_m^n (f_1^1)^{\otimes m}$$

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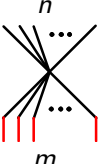
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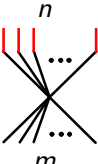
- As $KK_{n,m} \times 0 \leftrightarrow$  $\leftrightarrow \theta_m^n (f_1^1)^{\otimes m}$

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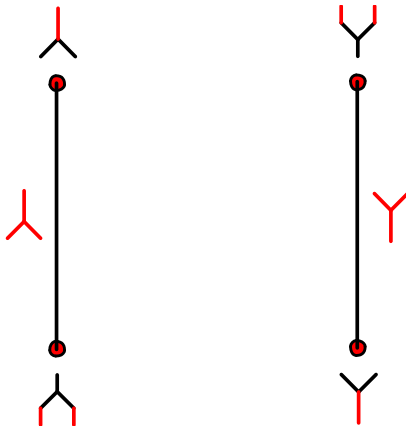
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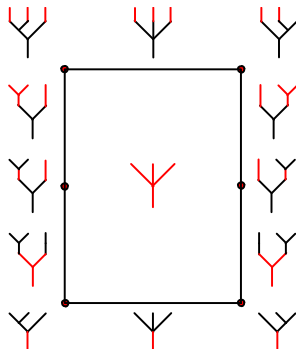
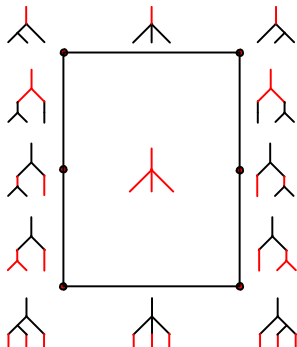
- As $KK_{n,m} \times 1 \leftrightarrow$  $\leftrightarrow (f_1^1)^{\otimes n} \theta_m^n$

- This latter cell plays an important role in our transfer algorithm

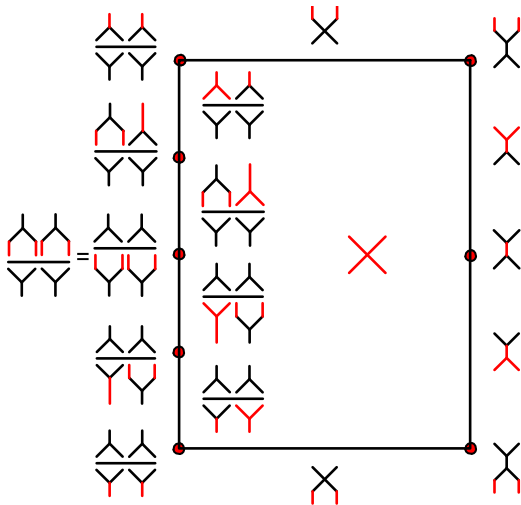
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- A morphism $A \Rightarrow B$ of A_∞ -bialgebras is the image

$$G = \{g_m^n \in \text{Hom}(A^{\otimes m}, B^{\otimes n})\}$$

of a map $\mathcal{J}\mathcal{J}_\infty \rightarrow U_{A,B}$ of relative matrads sending $f_m^n \mapsto g_m^n$

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- A morphism $G : A \Rightarrow B$ is an $\mathcal{H}_\infty$ -bimodule

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- ∇ -cocycles are chain maps $V \rightarrow W$
- A chain map $g : A \rightarrow B$ of DGMs induces a cochain map

$$\tilde{g} : \mathcal{E}nd_{TA} \rightarrow U_{A,B}$$

defined on $u \in \text{Hom}(A^{\otimes m}, A^{\otimes n})$ by $\tilde{g}(u) = g^{\otimes n} u$

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- **Transfer Theorem** *If $\tilde{g}$ is a homology isomorphism, then*
 - *g induces an A_∞ -bialgebra structure $\omega_A = \{\omega_A^{n,m}\}$ on A and*
 - *g extends to an A_∞ -bialgebra map $G = \{g_m^n \mid g_1^1 = g\} : A \Rightarrow B$*

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- A chain map/homology isomorphism $g : A \rightarrow B$ such that $\tilde{g}$ is a homology isomorphism

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- Define $\beta : C_0(JJ_{1,1}) \rightarrow \text{Hom}(A, B)$ by

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$$\theta_2^1(f_1^1 \otimes f_1^1) = \text{red vertex diagram} \mapsto \omega_B^{1,2}(g \otimes g)$$

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- Define $\alpha_A : C_0(\partial KK_{1,3}) \rightarrow Hom(A^{\otimes 3}, A)$ by

$$\text{Y} \mapsto \omega_A^{1,2}(\omega_A^{1,2} \otimes \mathbf{1}) \text{ and } \text{Y} \mapsto \omega_A^{1,2}(\mathbf{1} \otimes \omega_A^{1,2})$$

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- Extend β to the vertex $\text{Y} \subset JJ_{1,2}$ via $\text{Y} \mapsto g\omega_A^{1,2}$

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$$\text{Y} \mapsto (\omega_A^{2,1} \otimes \mathbf{1}) \omega_A^{2,1} \text{ and } \text{Y} \mapsto (\mathbf{1} \otimes \omega_A^{2,1}) \omega_A^{2,1}$$

- Extend β to the vertex $\text{Y} \subset JJ_{1,2}$ via $\text{Y} \mapsto (g \otimes g) \omega_A^{2,1}$

The Transfer Algorithm: Initialization

- Define $\alpha_A : C_0(\partial KK_{2,2}) \rightarrow \text{Hom}(A^{\otimes 2}, A^{\otimes 2})$ by

$$\text{Diagram} \mapsto \left(\omega_A^{1,2} \otimes \omega_A^{1,2} \right) \sigma_{2,2} \left(\omega_A^{2,1} \otimes \omega_A^{2,1} \right)$$

and

$$\text{Diagram} \mapsto \omega_A^{2,1} \omega_A^{1,2}$$

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- Note that $\left[\omega_B^{1,2} (g \otimes g) - g \omega_A^{1,2} \right] = 0$

The Transfer Algorithm: Initialization

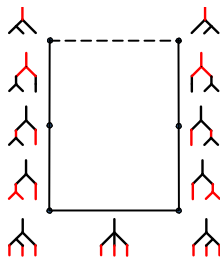
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- Define β on $\partial JJ_{1,3} \setminus \text{int } KK_{1,3} \times 1$:



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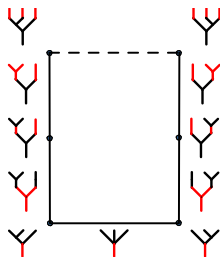
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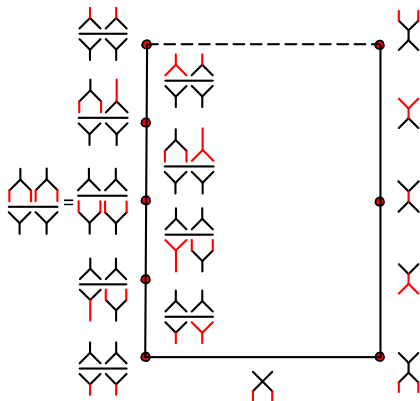
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The Transfer Algorithm: Initialization

Having defined β on $\text{red trivalent vertex}$ and red Y-vertex , define β on $\partial JJ_{2,2} \setminus \text{int } KK_{2,2} \times 1$:



The Transfer Algorithm: Induction Hypothesis

Given $m + n \geq 4$, assume for $i + j < m + n$, $ij \neq 1$, there exist maps

- $\alpha_A : C_*(KK_{j,i}) \rightarrow \text{Hom}(A^{\otimes i}, A^{\otimes j})$ of matrads sending $\theta_i^j \mapsto \omega_A^{j,i}$

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The Transfer Algorithm: Induction

For each $i + j = m + n$, $ij \neq 1$

- Define α_A on a monomial in $C_*(\partial KK_{n,m})$ to be its corresponding fraction product of operations in $\{\omega_A^{j,i}\}$; let

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- Define $\beta \left((f_1^1)^{\otimes n} \theta_m^n \right) = g^{\otimes n} \omega_A^{n,m}$

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- This completes the induction

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- $A = H^*(X; \mathbb{Z}_2) = \{1, a_2, a_3, b, a_2 a_3 = Sq^2 b, \dots\}$

A homotopy Gerstenhaber algebra structure on A

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- Induced map $\phi = E'_{1,0} + E'_{0,1} + E'_{1,1} : BA \otimes BA \rightarrow A$ acts nontrivially:

$$\phi([b] \otimes [b]) = a_2 a_3$$

Induced multiplication on BA

- ϕ lifts to a chain map $\mu : BA \otimes BA \rightarrow BA$ of DG coalgebras given by

$$\mu = \sum_{k \geq 0} \downarrow^{\otimes k+1} \phi^{\otimes k+1} \bar{\psi}^{(k)}$$

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- Denote: $\alpha_i = \text{cls}[a_i] \quad \beta = \text{cls}[b]$
- Choose a cycle-selecting map $g : H \rightarrow BA$ such that

$$g(\text{cls}[x_1 | \cdots | x_n]) = [x_1 | \cdots | x_n]$$

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Non-operadic operation defined

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$$\begin{aligned}\nabla g_2^2 &= (\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi \\ &\quad + (\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2 \\ &\quad + (g \otimes g) \omega_2^2\end{aligned}$$

- $((\mu \otimes \mu) \sigma_{2,2} (\Delta g \otimes g^2 + g^2 \otimes (g \otimes g) \Delta_H) + g^2 \varphi) (\beta \otimes \beta) = 0$
- $((\mu (g \otimes g) \otimes g_2 + g_2 \otimes g \varphi) \sigma_{2,2} (\Delta_H \otimes \Delta_H) + \Delta g_2) (\beta \otimes \beta)$
 $= [a_2] \otimes [a_3]$
- $[a_2] \otimes [a_3]$ represents a non-zero class in $BA \otimes BA$

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 $= [a_2] \otimes [a_3]$
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- Thus no choice of g_2^2 kills the obstruction $[a_2] \otimes [a_3]$

Conclusion:

- *There is an operation ω_2^2 and a cochain homotopy g_2^2 satisfying the relation on $JJ_{2,2}$ such that*

$$\omega_2^2(\beta \otimes \beta) = \alpha_2 \otimes \alpha_3$$

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- We have identified the first non-operadic cohomology operation

Thank you!